

Pricing Multi-Unit Markets

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We study the power and limitations of posted prices in multi-unit markets, where agents arrive sequentially in an arbitrary order. We prove upper and lower bounds on the largest fraction of the optimal social welfare that can be guaranteed with posted prices under a range of assumptions about the designer's information and agents' valuations. Our results provide insights about the relative power of uniform and non-uniform prices, the relative difficulty of different valuation classes, and the implications of different informational assumptions. Among other results, we prove constant-factor guarantees for agents with subadditive valuations over identical items, even in an incomplete-information setting and with uniform prices. We also show that no constant-factor guarantee is possible for general valuations over identical items, even in a full-information setting and with non-uniform prices.

CCS Concepts: • **Theory of computation** → **Algorithmic mechanism design; Computational pricing and auctions;**

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1 INTRODUCTION

We consider the problem of allocating identical items to agents with the objective of maximizing social welfare. There are a number of identical items, each agent has a valuation describing her value for a given number of items, and our goal is to compute nonnegative and integral quantities of items that should be assigned to each agent to maximize the total value over all agents.

This problem underlies the design of *multi-unit auctions*, which have played a starring role in the fields of classical and algorithmic mechanism design and in both theory and practice. As with any

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welfare-maximization problem, the problem can be solved in principle using the VCG mechanism. There has been extensive work on the design and analysis of more practical multi-unit auctions. There are indirect implementations of the VCG mechanism, most famously Ausubel’s ascending *clinching auction* for downward-sloping (a.k.a. submodular) valuations [Ausubel 2004]. Work in algorithmic mechanism design has identified mechanisms that retain the dominant-strategy incentive-compatibility of the VCG mechanism while running in time polynomial in the number of agents and the logarithm of the number of items (rather than polynomial in both the number of agents and items), at the cost of a bounded loss in the social welfare. Indeed, Nisan [2015] argues that the field of algorithmic mechanism design can be fruitfully viewed through the lens of multi-unit auctions.

The multi-unit auction formats used in practice typically sacrifice dominant-strategy incentive-compatibility in exchange for simplicity and equitability; a canonical example is the uniform-price auctions suggested by Milton Friedman (see Friedman [1991]) and used (for example) by the U.S. Treasury to sell government securities. Uniform-price auctions do not always maximize social welfare (e.g., because of demand reduction), but they do admit good “price-of-anarchy” guarantees [Markakis and Telelis 2015], meaning that every equilibrium results in social welfare close to the maximum possible.

A key drawback of all of the mechanisms above is that they require all agents to participate simultaneously to coordinate their allocations and respect the supply constraint. For example, in a uniform-price auction, all of the agents’ bids are used to compute a market-clearing price-per-unit, which then determines the allocations of all of the agents. It is evident from our daily experience that, in many different markets, buyers arrive and depart asynchronously over time, making purchasing decisions as a function of their preferences and the current prices of the goods for sale.¹ The goal of this article is to develop theory that explains the efficacy of such *posted prices* in markets where agents arrive sequentially rather than simultaneously and that gives guidance on how to set prices to achieve an approximately welfare-maximizing outcome.

1.1 The Model

We consider a setting where a designer must post prices in advance, before the arrival of any agents. We assume that the supply is known. The designer is given full or incomplete information about agents’ valuations, and must then set a price for each item.² Agents then arrive in an arbitrary (worst-case) order, with each agent taking a utility-maximizing bundle (breaking ties arbitrarily), given the set of items that remain. These prices are *static* in that they remain fixed throughout the entire process.

Example 1.1. Suppose there are three items and two agents, each with value 5 for getting one item, value 9 for getting two items, and value 11 for getting all three items, and suppose a designer prices every item at 4. The first agent will choose either one or two items (breaking the tie arbitrarily). If the first agent chooses 2 items, then the second agent will take the only item remaining; if the first agent chooses one item, then the second agent will take either one or two items.

In general, we allow different items to receive different prices (as will be the case in the VCG mechanism for this problem, for example.) With identical items, however, it is natural to focus on *uniform prices*, where every item is given the same price. Generally speaking, we are most interested in positive results for uniform prices and negative results for non-uniform prices.

¹For examples involving identical items, think about general-admission concert tickets, pizzas at Una Pizza Napoletana (which shuts down for the night when the dough runs out), or shares in an IPO (other than Google [Ritter 2014]).

²No non-trivial guarantees are possible without at least partial knowledge about agents’ valuations.

Table 1. Summary of Results

(a) Full information

	Uniform prices	Non-uniform prices
Submodular	$\frac{1}{2}$ (4.7, 4.8, 4.9)	$\frac{2}{3}$ (4.2, 4.3) [2 items] $\geq \frac{5}{7} - \frac{1}{m}$ (4.4), ≤ 0.802 (4.5) [m items]
XOS	$\frac{1}{2}$ (8.2)	$\leq 1 - \frac{1}{e}$ (5.1)
Subadditive	$\frac{1}{3}$ (6.1, 6.7) $\frac{2}{3}$ (6.5, 6.9) [2 identical buyers]	$\leq \frac{1}{2}$ (6.6) [even with 2 buyers] $\leq \frac{3}{4}$ (6.8) [even with 2 identical buyers]
General	$\frac{1}{m}$ (7.1)	$\frac{1}{m}$ (7.2)

(b) Incomplete information

	Uniform prices	Non-uniform prices
XOS	$\frac{1}{2}$ (8.2)	$\leq 1 - \frac{1}{e}$ (5.1)
Subadditive	$\geq \frac{1}{4}$ (8.4)	$\leq \frac{1}{2}$ (6.6) [even with 2 buyers] $\leq \frac{3}{4}$ (6.8) [even with 2 identical buyers]

All results are new to this article. Numbers in parentheses refer to the corresponding theorem or proposition number. The parameter m denotes the number of items.

The overarching goal of this article is to characterize the largest fraction of the optimal social welfare that can be guaranteed with posted prices, under a range of assumptions about the designer's information and agents' valuations. This goal is inherently quantitative, but our results also provide qualitative insights about the relative power of uniform and non-uniform prices, the relative difficulty of different valuation classes, and the implications of different informational assumptions. For instance, we show that with uniform prices in the full-information setting, sub-additive valuations are harder than XOS valuations, whereas XOS and submodular valuations are equally hard (see Section 2 for precise definitions of these valuation classes). Another example is that while non-uniform prices can be used to obtain a better guarantee for submodular valuations, this is not the case for general valuations.

1.2 Our Results

The majority of our results are summarized in Table 1; we highlight a subset of these next. First, consider the case of a Bayesian setting with XOS agent valuations (see Section 2 for definitions). That is, each agent's valuation is drawn independently from a known (possibly agent-specific) distribution over XOS valuations. Feldman et al. [2015] show that, even with non-identical items, posted prices can always obtain expected welfare at least $1/2$ times the maximum possible. This factor of $1/2$ is tight, even for the special case of a single item and i.i.d. agents. The posted prices used by Feldman et al. [2015] are non-uniform, even when the result is specialized to the case of identical items (the price of an item is based on its expected marginal contribution to an optimal allocation, which can vary across items). We prove in Theorem 8.2 that with identical items, and agents with independent (not necessarily identical) XOS valuations, uniform prices suffice to achieve the best-possible guarantee of half the optimal expected welfare. Moreover, this result extends to any class of valuations that is c -close to XOS valuations, with an additional loss of a factor of c (Theorem 8.3).

While the $1/2$ -approximation above is tight for an incomplete-information setting, this problem is already interesting in the full-information case where the buyers' valuations are known (with the order of arrival still worst-case). Can we improve over the approximation factor of $1/2$ under

this stronger informational assumption? We prove that uniform prices cannot achieve an approximation factor better than $1/2$, even for the more restrictive class of submodular valuations and even with two agents (Proposition 4.8) or identical agents (Proposition 4.9). In contrast, with non-uniform prices (still for submodular valuations), we prove that an approximation of $2/3$ is possible (Theorem 4.2), which is again tight.

We next consider the family of subadditive valuations, which strictly generalizes submodular and XOS valuations and is regarded as the most challenging class of valuations that forbid complements. For example, with non-identical items, it is not known whether or not posted prices can guarantee a constant fraction of the optimal social welfare. For identical items, we prove that this is indeed possible. In the incomplete-information setting (and identical items), we show that subadditive valuations are 2-close to XOS valuations (Proposition 3.4), which leads to approximation factor of $1/4$ (Theorem 8.4). We can also do better in the full-information setting: Uniform prices can guarantee a $1/3$ fraction of the optimal social welfare (Theorem 6.1), and the approximation is tight (Proposition 6.7), while even non-uniform prices cannot guarantee a factor bigger than $1/2$, even with only two agents (Proposition 6.6). In the case of two identical agents, uniform prices can guarantee a $2/3$ fraction of the optimal welfare (Theorem 6.5), and this is tight (Proposition 6.9).

With all these positive results, the reader might wonder whether constant factor guarantees can be provided for general valuations. Unfortunately, this is not the case. For general valuations, we show that even in the full-information setting and with non-uniform prices, and even when there are only two agents and the arrival order is known, posted prices cannot guarantee more than a $1/m$ fraction of the optimal social welfare (Proposition 7.1), where m denotes the number of items. If the seller can control the arrival order, however, then even uniform prices can guarantee half of the optimal social welfare (Theorem 7.3). No better bound is possible, even for identical valuations and with non-uniform prices (Proposition 7.4).

Besides the results themselves, a recurring technique that we develop in this article is that of considering the multiset containing the marginal values of all agents and setting appropriate prices based on this multiset. This is a powerful technique that, as we demonstrate throughout the article, can be used to obtain a broad range of guarantees in pricing multi-unit markets. For instance, it is often helpful to set prices just below or just above the threshold where enough items will be sold based on the marginal values. However, we will also need to carefully choose the number of items to be sold at each price, and this is at times a challenging task. For submodular valuations, we can consider the multiset of marginal values directly, since every agent has a decreasing marginal value for each subsequent item. However, for subadditive and general valuations, we will need to work with the “submodular closures” of these valuations, which correspond to the minimal submodular valuations that upper bound our valuations. We believe that these techniques can be useful for deriving results in general settings with identical items.

1.3 Further Related Work

The design and analysis of simple mechanisms has been an active area of study in algorithmic mechanism design, particularly within the last decade. This focus is motivated in part by the observation that simple mechanisms are highly desired in practical scenarios. Examples of simple mechanisms that are used in practice are the generalized second price auctions (GSP) for online advertising [Edelman et al. 2007; Lucier and Paes Leme 2011; Lucier et al. 2012; Paes Leme and Tardos 2010; Varian 2007] and simultaneous item auctions (where the agents bid separately and simultaneously on multiple items) [Bhawalkar and Roughgarden 2012; Christodoulou et al. 2008; Feldman et al. 2013; Hassidim et al. 2011]. These mechanisms are not truthful and are evaluated in equilibrium using the price of anarchy measure.

Posted price mechanisms constitute perhaps the most prevalent method for selling goods in practice. By simply publishing prices for individual items, these mechanisms are extremely easy

to understand and participate in. It should therefore not come as a surprise that posted price mechanisms have been studied extensively for various objective functions (e.g., welfare, revenue, makespan), information structures of values (e.g., full-information, Bayesian, online), and valuation classes (e.g., unit-demand, submodular, XOS). For example, a long line of work has focused on sequential posted prices for revenue maximization and has shown, among other things, that a form of posted price mechanisms can achieve a constant fraction of the optimal revenue for agents with unit-demand valuations [Chawla et al. 2007, 2010a, 2010b]. Revenue maximization with sequential posted prices has also been studied for a single item, both in large markets [Blumrosen and Holenstein 2008] and when the distributions are unknown [Babaiouff et al. 2011], for additive valuations [Babaiouff et al. 2014; Bateni et al. 2015], and for a buyer with complements [Eden et al. 2017]. Dütting et al. [2017] provide a general framework for posted price mechanisms. In several of these works, posted price mechanisms are allowed to discriminate between agents and set different prices for each of them. In contrast, in this work we do not consider discriminatory prices.

Another line of research relevant to our work considers market equilibria, for example, those achieved by Walrasian prices. A result of Kelso Jr. and Crawford [1982] states that for the class of gross-substitute valuations, there always exists a Walrasian equilibrium, meaning that one can assign prices to items so as to achieve the optimal social welfare. However, this result is based on the assumption that agents break ties in a particular way. As such, the existence of Walrasian prices does not carry over welfare guarantees to our setting, even for unit-demand valuations. We believe that the worst-case perspective that we take is more realistic in our setting, where we do not have control over how agents break ties.

In addition to the aforementioned works, a new line of research has considered dynamic posted prices in online settings such as for the k -server and parking problems [Cohen et al. 2015]. Moreover, posted price mechanisms have been studied in the context of welfare maximization in matching markets, where prices are dynamic (i.e., can change over the course of the mechanism) but do not depend on the identity of the agents [Cohen-Addad et al. 2016]. With static prices, it was recently shown that one can achieve strictly more than half of the welfare in the full information setting with binary unit-demand valuations [Eden et al. 2018].

The sequential arrival of agents considered in posted price mechanisms fits into the framework of online mechanisms, which deals with dynamic environments with multiple agents having private information [Babaiouff et al. 2015; Hajiaghayi et al. 2004; Parkes 2007]. Our work shows that for identical items and agents with subadditive valuations, posted prices can guarantee a constant fraction of the welfare even while setting the (uniform) prices up front.

2 PRELIMINARIES

We consider a setting with a set M of m identical items and a set N of n buyers. Each buyer has a valuation $v_i : 2^M \rightarrow \mathbb{R}_{\geq 0}$ that indicates his value for every set of items. Since items are identical, the valuation depends only on the number of items. We assume that valuations are non-decreasing (i.e., $v_i(T) \leq v_i(S)$ for $T \subseteq S$) and normalized (i.e., $v_i(\emptyset) = 0$). We use $v_i(S|T) = v_i(S \cup T) - v_i(T)$ to denote the marginal value of bundle S given bundle T .

A buyer valuation profile is denoted by $\mathbf{v} = (v_1, \dots, v_n)$. An allocation is a vector of disjoint sets $\mathbf{x} = (x_1, \dots, x_n)$, where x_i denotes the bundle associated with buyer $i \in [n]$ (note that it is not required that all items are allocated). As with valuations, since we consider identical items, an allocation can be represented by the number of items allocated to each buyer. The social welfare (SW) of an allocation \mathbf{x} is $SW(\mathbf{x}, \mathbf{v}) = \sum_{i=1}^n v_i(x_i)$, and the optimal social welfare is denoted by $OPT(\mathbf{v})$. When clear from the context we omit \mathbf{v} and write SW and OPT for the social welfare and optimal social welfare, respectively.

For two valuations v, v' , we say that $v \geq v'$ iff $v(S) \geq v'(S)$ for every set S . A hierarchy over complement-free valuations is given by Lehmann et al. [2006].

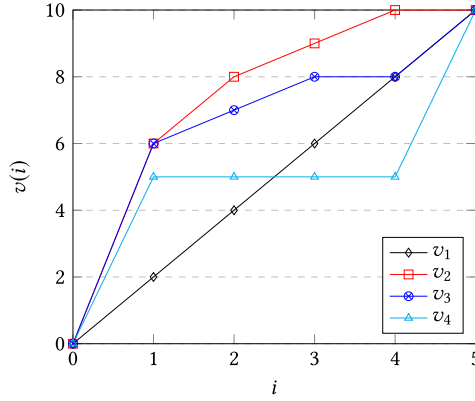


Fig. 1. Examples of valuations belonging to the different classes in Definition 2.3. Valuation v_1 is additive, v_2 is submodular, v_3 is XOS, and v_4 is subadditive. Each valuation besides v_1 does not belong to any of the classes below its class.

Definition 2.1. A valuation v is

- *additive* if $v(S) = \sum_{i \in S} v(\{i\})$ for every set $S \subseteq M$.
- *submodular* if $v(\{i\}|S) \geq v(\{i\}|T)$ for every item $i \notin T$ and sets S, T such that $S \subseteq T \subseteq M$.
- *XOS* if there exist additive valuations v^1, \dots, v^k such that $v(S) = \max_{j=1, \dots, k} v^j(S)$ for every set $S \subseteq M$.
- *subadditive* if $v(S) + v(T) \geq v(S \cup T)$ for any sets $S, T \subseteq M$.

Since we assume throughout the article that all items are identical, we only work with valuations over identical items.

Definition 2.2. A valuation v is said to be a *valuation over identical items* if $v(S) = v(T)$ for every sets $S, T \subseteq M$ such that $|S| = |T|$. A valuation over identical items can thus be represented by a non-decreasing valuation $v : \{0, 1, \dots, m\} \rightarrow \mathbb{R}_{\geq 0}$, which assigns a non-negative real value to any integer in $[m]$ (recall $v(0) = 0$ as we assume normalized valuations).

Unless specified otherwise, all valuations in this article are assumed to be valuations over identical items.

In what follows, we adjust the definitions of additive, submodular, XOS, and subadditive valuations in Definition 2.1 to the case of valuations over identical items. The simplified definition for XOS valuations follows from the equivalence between XOS and fractional subadditivity [Feige 2009].

Definition 2.3. A valuation over identical items v is said to be

- *additive* if $v(i) = a \cdot i$ for every integer $0 \leq i \leq m$ for some constant a .
- *submodular* if $v(i) - v(i-1) \geq v(i+1) - v(i)$ for every integer $1 \leq i \leq m-1$.
- *XOS* if $v(i) \geq \frac{i}{j} \cdot v(j)$ for any integers $1 \leq i < j \leq m$.
- *subadditive* if $v(i) + v(j) \geq v(i+j)$ for any integers $1 \leq i, j \leq m$ with $i+j \leq m$.

Examples of valuations belonging to these classes are shown in Figure 1.

We assume that the agents arrive sequentially. We will for the most part set static prices for the items, and each arriving agent takes a bundle from the remaining items that maximizes her utility, with ties broken arbitrarily. For some results, we will assume dynamic prices, i.e., the seller can set new prices for the remaining items for each iteration (but without knowing which agent will

arrive next). If prices $\mathbf{p} = (p_1, \dots, p_m)$ are set on the m items, and an agent buys a subset S of them, then her utility is given by $v(|S|) - \sum_{i \in S} p_i$. For most of the article, we will assume that the arrival order of the agents is unknown, but we will also consider settings where we know this order or where we even have control over the order. We are interested in the social welfare that we can obtain by setting prices in comparison to the optimal social welfare with respect to the worst case arrival order.

Given a buyer valuation profile with submodular valuations v_1, \dots, v_n , we will often consider the multiset $V = \{v_i(j) - v_i(j-1) \mid i \in [n], j \in [m]\}$, which consists of all marginal values for all buyers and items. Let δ be the minimum positive difference between any two values in V , and let $\epsilon = \delta/2$. Note that the optimal allocation is to sort the marginal values in V in non-increasing order, and allocate items in that order. Thus, OPT is the sum of the m highest marginal values in V .

For every value x , let $G(x)$ be the number of marginal values in V that are strictly greater than x , and let $E(x)$ be the number of marginal values in V that are equal to x . Without loss of generality, we will assume that $G(0) \geq m$, since otherwise we can obtain the optimal social welfare by setting all prices to ϵ , which guarantees that all $G(0)$ items yielding positive marginal utility are sold. Let b be the unique value satisfying $G(b) < m$ and $G(b) + E(b) \geq m$. Let $m' = G(b)$; i.e., m' is the number of marginal values strictly greater than b .

Example 2.1. Suppose $m = 3$ and there are two agents. The first agent has the valuation $v(1) = 5$, $v(2) = 9$, and $v(3) = 11$, while the second agent has the valuation $v(1) = 2$, $v(2) = 4$, and $v(3) = 5$. Then we have $V = \{5, 4, 2, 2, 2, 1\}$, $\delta = 1$, and $\epsilon = 0.5$. The optimal welfare is 11 (obtained by either allocating all three items to the first agent, or by allocating two items to the first agent and a single item to the second agent). Moreover, we have $G(2) = 2$, $E(2) = 3$, $b = 2$, and $m' = G(2) = 2$.

3 PROPERTIES OF VALUATIONS OVER IDENTICAL ITEMS

In this section, we consider properties of valuations over identical items. In addition to being interesting in their own right, these properties will later help us establish welfare guarantees for posted prices (Theorem 8.4).

First, we show that every valuation (over identical items) admits a unique minimal XOS valuation as well as a unique minimal submodular valuation that upper bounds it. We will sometimes refer to the latter valuations as the ‘‘closures’’ of the former valuation.

PROPOSITION 3.1. *Let v be a valuation over a bounded domain, let \tilde{v} be the valuation $\tilde{v}(i) = \max_{j \geq i} \frac{i}{j} \cdot v(j)$. Then \tilde{v} is XOS, $\tilde{v} \geq v$, and for every XOS valuation \hat{v} such that $\hat{v} \geq v$, we have $\hat{v} \geq \tilde{v}$.*

PROOF. We start by proving that \tilde{v} is an XOS valuation. By Definition 2.3 it is enough to show that for every i, j such that $i < j$ it holds that $\tilde{v}(i) \geq \frac{i}{j} \cdot \tilde{v}(j)$. Let i, j be such that $i < j$. We know that there exists $k \geq j$ such that $\tilde{v}(j) = \frac{j}{k} \cdot v(k)$ by the definition of \tilde{v} , and therefore $\tilde{v}(i) \geq \frac{i}{k} \cdot v(k) = \frac{i}{k} \cdot \frac{k}{j} \cdot \tilde{v}(j) = \frac{i}{j} \cdot \tilde{v}(j)$, as needed. It holds that $\tilde{v} \geq v$, since $\tilde{v}(i) \geq \frac{i}{i} \cdot v(i)$ for all i .

Now, let \hat{v} be an XOS valuation such that $\hat{v} \geq v$, and let $i \leq m$. Let $j = \operatorname{argmax}_{j \geq i} \frac{i}{j} \cdot v(j)$. We have that $\hat{v}(j) \geq v(j)$, and by Definition 2.3 that $\hat{v}(i) \geq \frac{i}{j} \cdot \hat{v}(j) \geq \frac{i}{j} \cdot v(j) = \tilde{v}(i)$, as needed. \square

PROPOSITION 3.2. *Let v be a valuation over a bounded domain, and let \tilde{v} be the valuation $\tilde{v}(i) = \max_{j \geq i, k \leq i} (\frac{i-k}{j-k} \cdot (v(j) - v(k)) + v(k))$ (when $j = i = k$ the expression in the maximum is $v(i)$). Then \tilde{v} is submodular, $\tilde{v} \geq v$, and for every submodular valuation \hat{v} such that $\hat{v} \geq v$, we have $\hat{v} \geq \tilde{v}$.*

PROOF. We start by proving that \tilde{v} is a submodular valuation. To this end, we present an alternative way to define \tilde{v} that makes its submodularity clear. Plot the points $(i, v(i))$ for all $i \in [m]$ on a graph, and define $v'(0) = 0$. If we have defined $v'(j)$ for all $j \leq i$, then consider the least $k > i$

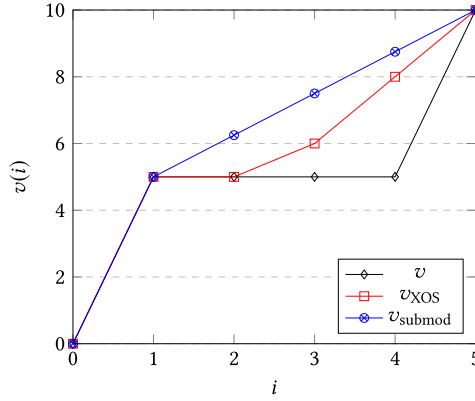


Fig. 2. An example of some valuation v and its XOS and submodular closures, as constructed in Propositions 3.1 and 3.2. The XOS closure of v is depicted by v_{XOS} , and its submodular closure is depicted as v_{submod} .

such that no point of v lies strictly above the line connecting $(i, v(i))$ and $(k, v(k))$. For all j with $i < j \leq k$, define $v'(j)$ so that the point $(j, v'(j))$ lies on this line. Since the slope of the lines cannot increase along this process, v' is submodular.

It holds that $\tilde{v} \geq v$, since $\tilde{v}(i) \geq v(i)$ for all i . We show that in fact $\tilde{v} = v'$. Fix any $i \in [m]$, and suppose that $v'(i)$ is defined by the line connecting $(r, v(r))$ and $(s, v(s))$ for some $r < s$. We have $v'(i) = v(r) + (v(s) - v(r)) \cdot \frac{i-r}{s-r}$, so $v'(i) \leq \tilde{v}(i)$. Moreover, since all points of v lie on or below this line, the maximum of $v(j) + (v(k) - v(j)) \cdot \frac{i-j}{k-j}$ is attained at $j = r$ and $k = s$, so $v'(i) = \tilde{v}(i)$.

Now, let \hat{v} be any submodular valuation such that $\hat{v} \geq v$, and let $i \leq m$. Let $r, s = \operatorname{argmax}_{j \leq i, k \geq i} v(j) + (v(k) - v(j)) \cdot \frac{i-j}{k-j}$. We have that $\hat{v}(r) \geq v(r)$ and $\hat{v}(s) \geq v(s)$. Therefore,

$$\begin{aligned} \hat{v}(i) &\geq \hat{v}(r) + (\hat{v}(s) - \hat{v}(r)) \cdot \frac{i-r}{s-r} = \hat{v}(r) \cdot \frac{s-i}{s-r} + \hat{v}(s) \cdot \frac{i-r}{s-r} \\ &\geq v(r) \cdot \frac{s-i}{s-r} + v(s) \cdot \frac{i-r}{s-r} = v(r) + (v(s) - v(r)) \cdot \frac{i-r}{s-r} = \tilde{v}(i), \end{aligned}$$

as needed. \square

Examples of a valuation along with their XOS and submodular closures can be found in Figure 2.

We are interested in approximating valuations with “simpler” valuations. Specifically, for two classes of valuations $\mathcal{V}_1 \subseteq \mathcal{V}_2$, we want to determine the smallest constant c such that for any valuation $v \in \mathcal{V}_2$, there exists a valuation $\tilde{v} \in \mathcal{V}_1$ such that $v \leq \tilde{v} \leq cv$. We answer this question for each pair from the classes of subadditive, XOS, and submodular valuations and show that the best constant is $c = 2$ for all of these pairs. (Note that since all three classes are closed under scalar multiplication, the inequality $v \leq \tilde{v} \leq cv$ above can also be replaced by $v/c \leq \tilde{v} \leq v$.)

PROPOSITION 3.3. *For every subadditive valuation v , there exists a submodular valuation \tilde{v} such that $\tilde{v} \geq v$ and $\tilde{v} \leq 2 \cdot v$.*

PROOF. Consider the minimal submodular valuation \tilde{v} that upper bounds v defined in Proposition 3.2. For a certain i let $s, r = \operatorname{argmax}_{j \geq i, k \leq i} \frac{i-k}{j-k} \cdot (v(j) - v(k)) + v(k)$. It holds that

$$\begin{aligned} \tilde{v}(i) &= \frac{i-r}{s-r} \cdot (v(s) - v(r)) + v(r) \\ &\leq \frac{i}{s} \cdot (v(s) - v(r)) + v(r) \end{aligned}$$

$$\begin{aligned}
&\leq \frac{i}{s} \cdot (v(i \cdot \lceil s/i \rceil) - v(r)) + v(r) \\
&\leq \frac{i}{s} \cdot (\lceil s/i \rceil \cdot v(i) - v(r)) + v(r) \\
&\leq \frac{i}{s} \left(\left(\frac{s}{i} + 1 \right) \cdot v(i) - v(r) \right) + v(r) \\
&= \frac{i}{s} \left(\left(\frac{s}{i} + 1 \right) \cdot v(i) \right) + v(r) \left(1 - \frac{i}{s} \right) \\
&\leq \frac{i}{s} \left(\left(\frac{s}{i} + 1 \right) \cdot v(i) \right) + v(i) \left(1 - \frac{i}{s} \right) \\
&= 2 \cdot v(i),
\end{aligned}$$

where the second and the last inequalities follow from the fact that v is nondecreasing, and the third inequality follows from the fact that v is subadditive. \square

Proposition 3.3 immediately yields the following two results.

PROPOSITION 3.4. *For every subadditive valuation v , there exists an XOS valuation \tilde{v} such that $\tilde{v} \geq v$ and $\tilde{v} \leq 2 \cdot v$.*

PROPOSITION 3.5. *For every XOS valuation v , there exists a submodular valuation \tilde{v} such that $\tilde{v} \geq v$ and $\tilde{v} \leq 2 \cdot v$.*

Next, we show that the constant 2 is the best possible for all three pairs of classes. Proposition 3.7 follows directly from Proposition 3.6.

PROPOSITION 3.6. *For every $0 < c < 2$, there exists an XOS valuation v such that for every submodular valuation \tilde{v} with $v \leq \tilde{v}$, we have $v(i) < \tilde{v}(i)/c$ for some i .*

PROOF. Let l be a large integer to be chosen later. Consider the XOS valuation

$$v(i) = \begin{cases} 1 & \text{if } i \in \{1, \dots, \ell\} \\ \frac{i}{\ell} & \text{if } i \in \{\ell, \dots, \ell^2\}. \end{cases}$$

By Proposition 3.2 and setting $j = 1, k = \ell^2$, we have that the minimal submodular valuation \tilde{v} that upper bounds v satisfies

$$\tilde{v}(\ell) \geq v(1) + \frac{\ell - 1}{\ell^2 - 1} \cdot (v(\ell^2) - v(1)) = 1 + \frac{\ell - 1}{\ell + 1}.$$

Setting ℓ large enough so that $\frac{\ell - 1}{\ell + 1} > c - 1$ yields the desired result. \square

PROPOSITION 3.7. *For every $0 \leq c < 2$, there exists a subadditive valuation v such that for every submodular valuation \tilde{v} with $v \leq \tilde{v}$, we have $v(i) < \tilde{v}(i)/c$ for some i .*

PROPOSITION 3.8. *For every $0 \leq c < 2$, there exists a subadditive valuation v such that for every XOS valuation \tilde{v} with $v \leq \tilde{v}$, we have $v(i) < \tilde{v}(i)/c$ for some i .*

PROOF. Let l be a large integer to be chosen later. Consider the subadditive valuation

$$v(i) = \begin{cases} 1 & \text{if } i \in \{1, \dots, l\} \\ 2 & \text{if } i = l + 1 \end{cases}.$$

By Proposition 3.1, we have that the minimal XOS valuation \tilde{v} that upper bounds v satisfies $\tilde{v}(l) = \frac{2l}{l+1}$. Setting l high enough so that $\frac{2l}{l+1} > c$ yields the desired result. \square

4 SUBMODULAR VALUATIONS

In this section, we consider submodular valuations. We first show a reduction from submodular to unit-demand valuations. In particular, we show how to transform n submodular valuations into nm unit-demand valuations in a way that every static pricing for the m items that guarantees a fraction ρ of the optimal social welfare for the unit-demand valuations, will also guarantee a fraction ρ of the optimal social welfare for the submodular valuations. Given this reduction, it is without loss of generality to assume for positive results in Sections 4.1 and 4.2 that all valuations are unit-demand.

Let \mathcal{M} be a market with m items and n agents each with a valuation v_i . Let $OPT(\mathcal{M})$ be the social welfare of the welfare-maximizing allocation. Also, given prices \mathbf{p} , let $SW(\mathcal{M}, \mathbf{p})$ be the worst case social welfare obtained in market \mathcal{M} under pricing \mathbf{p} .

LEMMA 4.1. *Let \mathcal{M} be a market with n agents with submodular valuations v_i over $[m]$. Consider a market \mathcal{M}' with $n \cdot m$ agents, denoted by (i, j) for $i \in [n], j \in [m]$, with unit-demand valuations v_i^j with values $v_i(j) - v_i(j - 1)$. The following holds:*

- $OPT(\mathcal{M}) = OPT(\mathcal{M}')$
- For every pricing \mathbf{p} , $SW(\mathcal{M}, \mathbf{p}) \geq SW(\mathcal{M}', \mathbf{p})$.

PROOF. We prove the two parts in order.

- Consider an allocation of \mathcal{M} that allocates x_i items to agent i . The allocation in \mathcal{M}' that allocates items to agents (i, j) if $j \leq x_i$ yields the same social welfare, and therefore $OPT(\mathcal{M}) \leq OPT(\mathcal{M}')$. Conversely, consider an allocation in \mathcal{M}' that allocates items to a set of agents $A = \{(i_1, j_1), \dots, (i_m, j_m)\}$. The allocation in \mathcal{M} that allocates $\lfloor \sum_{j: (i, j) \in A} 1 \rfloor$ items to agent i yields at least the same social welfare, therefore $OPT(\mathcal{M}) \geq OPT(\mathcal{M}')$. We therefore have $OPT(\mathcal{M}) = OPT(\mathcal{M}')$.
- Let σ be the worst case order for market \mathcal{M} with pricing \mathbf{p} , where agent i arrives at turn $\sigma(i)$, and let (x_1, \dots, x_n) be the allocation that generates the worst-case social welfare. Let σ' be the order of arrival in market \mathcal{M}' where agent (i, j) arrives at turn $m \cdot (\sigma(i) - 1) + j$. One can check that the allocation where agent (i, j) is allocated an item if and only if $x_i \geq j$, is a valid allocation with respect to the order of arrival σ' and pricing \mathbf{p} . The social welfare of this allocation is $SW(\mathcal{M}, \mathbf{p})$. Hence $SW(\mathcal{M}, \mathbf{p}) \geq SW(\mathcal{M}', \mathbf{p})$.

This completes the proof. □

Remark. If there are no ties in the marginal values across all buyers (in fact, it is sufficient that the m th highest marginal value is different than the $(m + 1)$ -st highest marginal value), then the optimal social welfare can be achieved by setting all prices to be $b - \epsilon$, where b and ϵ are defined as in Section 2. However, as we will see (e.g., in Proposition 4.3), marginal value ties can cause the welfare from posted prices to drop significantly.

4.1 Non-uniform Pricing

We first show that we can obtain $2/3$ of the optimal welfare for submodular valuations if we are allowed to set non-uniform prices.

THEOREM 4.2. *For every market with submodular valuations, there exists a static item pricing \mathbf{p} that guarantees at least $2/3$ of the optimal social welfare.*

PROOF. By Lemma 4.1 it suffices to show the result for buyers with unit-demand valuations. Define b , m' , and ϵ as in Section 2. Consider the following pricing structures:

- (P1) set a price of $b - \epsilon$ to all items;
- (P2) set a price of $b - \epsilon$ to $m - m'$ items and a price of $b + \epsilon$ to m' items.

We claim that the maximum of these pricing structures gives at least $2/3$ of OPT .

Under pricing (P1) all m items will be sold, and each one has a marginal value of at least b ; thus, a social welfare of at least mb is guaranteed.

We next show that pricing (P2) guarantees a social welfare of at least $\max\{OPT - (m - m')b, OPT - m'b\}$.

The first observation is that every buyer whose value is greater than b will buy an item. This is because when the buyer arrives, at least one item with price $b + \epsilon$ is available (there are m' such items initially, and the other buyers can buy at most $m' - 1$ of them), and an item is worth to her more than $b + \epsilon$. Therefore, the social welfare of these items is at least $OPT - (m - m')b$. This also means that the social welfare under pricing (P2) is at least $OPT - (m - m')b$.

The second observation is that at least $m - 2m'$ items that generate value exactly b will be sold. (It is possible that $m - 2m'$ is negative, but this does not affect our analysis.) This follows from the fact that all $m - m'$ items with price $b - \epsilon$ will be sold (since these are the cheapest items and the demand is at least m), and at most m' of the $m - m'$ items will generate marginal value greater than b when sold. Therefore, the social welfare from these items is at least $(m - 2m')b$.

Since the sets of items considered in the first and second observations are disjoint, the total social welfare obtained by (P2) is at least their sum, which is $OPT - (m - m')b + (m - 2m')b = OPT - m'b$. So the social welfare under pricing (P2) is at least $\max\{OPT - (m - m')b, OPT - m'b\}$, as stated.

The average of the guarantee of pricing (P1) and the two guarantees of pricing (P2) is then $(mb + OPT - (m - m')b + OPT - m'b)/3 = 2OPT/3$. Therefore, the best of these two pricing structures gives social welfare at least $2OPT/3$. \square

The following proposition shows that this bound is tight.

PROPOSITION 4.3. *There exists a market with two items and two buyers with submodular valuations such that every static pricing can guarantee a social welfare of at most $2/3$ of the optimal social welfare.*

PROOF. Consider a market with two items and two buyers with the following valuations: Buyer 1 has a unit-demand valuation with a value of 2 for each item. Buyer 2 has an additive valuation with a value of 1 for each item. The optimal welfare is $OPT = 3$, obtained by allocating to each buyer a single item. We show that no pricing can guarantee more than a welfare of 2. Let $\mathbf{p} = (p_1, p_2)$ be some pricing. Distinguish between the following cases:

- (1) If $p_1, p_2 \leq 1$, then if the additive buyer arrives first, she will buy both items, resulting in a social welfare of 2.
- (2) If one of the prices is ≤ 1 and the other price is > 1 , then if the unit-demand buyer arrives first, she will buy the cheaper item. Then the additive buyer will not buy anything, resulting in a social welfare of 2.
- (3) If $p_1, p_2 > 1$, then the additive buyer will not buy any item, thus the social welfare cannot exceed 2.

This concludes the argument. \square

The negative result in Proposition 4.3 is obtained for a market with two items. In what follows, we show that the guaranteed social welfare is higher when the number of items is large.

THEOREM 4.4. *For every market of m items with submodular valuations, there exists a static item pricing \mathbf{p} that guarantees at least $5/7 - 1/m$ of the optimal social welfare.*

PROOF. By Lemma 4.1 it suffices to show the result for buyers with unit-demand valuations. Define V , b , m' , and ϵ as in Section 2. In addition, let $k = |\{x \in V \mid x \geq 2b\}|$. Recall that OPT is the sum of the m highest values in V , and therefore $\text{OPT} \geq mb$.

Let $f(i, V)$ be the function that returns the i th highest value in the multiset V . Let α_1 be the sum of the highest $\lfloor m'/2 \rfloor$ values in V (i.e., $\alpha_1 = \sum_{i \leq \lfloor m'/2 \rfloor} f(i, V)$), and let α_2 be the sum of the remaining $\lceil m'/2 \rceil$ values (i.e., $\alpha_2 = \sum_{\lfloor m'/2 \rfloor < i \leq m'} f(i, V)$). Consider the following pricing structures:

- (P1) set a price of $b - \epsilon$ to all items;
- (P2) set a price of $b - \epsilon$ to $m - m'$ items and a price of $b + \epsilon$ to m' items;
- (P3) set a price of $b - \epsilon$ to $m - k$ items and a price of $2b - \epsilon$ to k items;
- (P4) set a price of $b - \epsilon$ to $m - \lfloor m'/2 \rfloor$ items and a price of $b + \epsilon$ to $\lceil m'/2 \rceil$ items.

We claim that the maximum of these pricing structures gives at least $5/7 - 1/m$ of OPT .

Following the same reasoning as in the proof of Theorem 4.2, pricing (P1) guarantees a social welfare of at least mb and pricing (P2) guarantees a social welfare of at least $\max\{\text{OPT} - (m - m')b, \text{OPT} - m'b\}$. We next consider the new pricings (P3) and (P4).

We first show that pricing (P3) guarantees a social welfare of at least $\alpha_1 + (m - m')b$. Under pricing (P3) all k items with value at least $2b$ will be sold, and following the analysis in Theorem 4.2, at least $m - 2k$ additional items will be sold for a value of at least b . Thus, the social welfare is at least $\sum_{i \leq k} f(i, V) + (m - 2k)b$. We show that this expression is at least $\alpha_1 + (m - m')b$. We consider two cases:

- (1) $k < \lfloor m'/2 \rfloor$. In this case, we have:

$$\begin{aligned} & \sum_{i \leq k} f(i, V) + (m - 2k)b \\ &= \sum_{i \leq \lfloor m'/2 \rfloor} f(i, V) - \sum_{k < i \leq \lfloor m'/2 \rfloor} f(i, V) + 2b(\lfloor m'/2 \rfloor - k) + (m - 2\lfloor m'/2 \rfloor)b \\ &\geq \alpha_1 - \sum_{k < i \leq \lfloor m'/2 \rfloor} (f(i, V) - 2b) + (m - m')b \geq \alpha_1 + (m - m')b, \end{aligned}$$

where the last inequality follows from the definition of k (which implies that $f(i, V) < 2b$ for every $i > k$), and the remaining inequalities follow from simple algebraic manipulations.

- (2) $k \geq \lfloor m'/2 \rfloor$. In this case, we have:

$$\begin{aligned} & \sum_{i \leq k} f(i, V) + (m - 2k)b \\ &= \sum_{i \leq \lfloor m'/2 \rfloor} f(i, V) + \sum_{\lfloor m'/2 \rfloor < i \leq k} f(i, V) - 2b(k - \lfloor m'/2 \rfloor) + (m - 2\lfloor m'/2 \rfloor)b \\ &\geq \alpha_1 + \sum_{\lfloor m'/2 \rfloor < i \leq k} (f(i, V) - 2b) + (m - m')b \geq \alpha_1 + (m - m')b, \end{aligned}$$

where the last inequality follows from the definition of k (which implies that $f(i, V) \geq 2b$ for every $i \leq k$), and the remaining inequalities follow from simple algebraic manipulations.

Consider next pricing (P4). Distinguish between two cases:

- (1) All items are sold. In the worst case the following happens:

- (a) Items with price $b + \epsilon$ (there are $\lceil m'/2 \rceil$ such items) will generate social welfare from the lowest values in V that are at least $b + \epsilon$. Thus, the social welfare is at least $\sum_{\lceil m'/2 \rceil < i \leq m'} f(i, V) = \alpha_2$.
- (b) Items with price $b - \epsilon$ will generate social welfare of b each. Thus, the social welfare is at least $(m - \lceil m'/2 \rceil)b$.

Together, if all items are sold, then the social welfare is at least $\alpha_2 + (m - \lceil m'/2 \rceil)b$.

- (2) There exists an item that remains unsold. It follows that all values strictly greater than b contribute to the social welfare (otherwise, there would be no unsold item), yielding a contribution of at least $\alpha_1 + \alpha_2$ to the social welfare. In addition, all items of price $b - \epsilon$ must be sold. There are at least $m - m' - \lceil m'/2 \rceil$ such items beyond the ones considered toward $\alpha_1 + \alpha_2$. Together, if there exists an unsold item, then the social welfare is at least $\alpha_1 + \alpha_2 + (m - m' - \lceil m'/2 \rceil)b$.

Pricing (P4) guarantees the minimum of the two cases, which is

$$\min(\alpha_1 + \alpha_2 + (m - m' - \lceil m'/2 \rceil)b, \alpha_2 + (m - \lceil m'/2 \rceil)b).$$

Note that the first term is chosen from the minimum exactly when $\alpha_1 < m'b$.

We now show that the best of these four pricing schemes guarantees a social welfare of at least $(5/7 - 1/m) \cdot OPT$. Assume toward contradiction that all of the above pricing schemes obtain less than $(5/7 - 1/m) \cdot OPT$.

- From pricing (P2), we get that

$$(m - m')b \geq OPT - SW(P2) > \frac{2 \cdot OPT}{7}. \quad (1)$$

$$m'b \geq OPT - SW(P2) > \frac{2 \cdot OPT}{7}. \quad (2)$$

- From pricing (P3), we get that

$$\alpha_2 \geq OPT - SW(P3) > \frac{2 \cdot OPT}{7}. \quad (3)$$

We consider two cases:

- (1) $\alpha_1 > m'b$: In this case, we get from pricing (P4) that:

$$\begin{aligned} \alpha_1 - \lceil m'/2 \rceil b &> OPT - SW(P4) > \frac{2 \cdot OPT}{7} \implies \\ \alpha_1 &> \lceil m'/2 \rceil b + \frac{2 \cdot OPT}{7} \geq \frac{m'b}{2} - \frac{b}{2} + \frac{2 \cdot OPT}{7} \\ &\stackrel{(2)}{>} \frac{OPT}{7} - \frac{b}{2} + \frac{2 \cdot OPT}{7} = \frac{3 \cdot OPT}{7} - \frac{b}{2} \stackrel{(3)}{\implies} \\ \alpha_1 + \alpha_2 &\geq \frac{3 \cdot OPT}{7} - \frac{b}{2} + \frac{2 \cdot OPT}{7} = \frac{5 \cdot OPT}{7} - \frac{b}{2}. \end{aligned}$$

Since pricing (P2) guarantees a social welfare of $\alpha_1 + \alpha_2$, the last inequality shows that pricing (P2) obtains a social welfare of at least $(5/7 - 1/m) \cdot OPT$ (recall that $b \leq OPT/m$), a contradiction.

- (2) $\alpha_1 \leq m'b$: In this case, we get from pricing (P4) that:

$$\begin{aligned} \lceil m'/2 \rceil b &\geq OPT - SW(P4) > \frac{2 \cdot OPT}{7} \implies m'b \geq 2\lceil m'/2 \rceil b - b \geq \frac{4 \cdot OPT}{7} - b \implies \\ mb &= (m - m')b + m'b \stackrel{(1)}{\geq} \frac{2 \cdot OPT}{7} - b + \frac{4 \cdot OPT}{7} = \frac{6 \cdot OPT}{7} - b. \end{aligned}$$

Since pricing (P1) guarantees a social welfare of at least mb , the last inequality shows that pricing (P1) obtains a social welfare of at least $(5/7 - 1/m) \cdot OPT$ (recall that $b \leq OPT/m$), a contradiction.

This concludes the argument. \square

The guarantee in Theorem 4.4 approaches $5/7 \approx 0.714$ as the number of items grows. The next theorem shows that this bound cannot exceed 0.802 even for an arbitrarily large number of items.

THEOREM 4.5. *For every constant c , there exists a market with $m > c$ items with submodular valuations such that for any static item pricing \mathbf{p} , the social welfare guaranteed by the pricing is at most 0.802 of the optimal social welfare.*

PROOF. Consider a market with m items for large-enough m . We set $\beta \approx 2.24698$ to be a solution to the equation $x^3 - 2x^2 - x + 1 = 0$, and $\alpha = \frac{1}{(\beta-1)^2} \approx 0.643104$. An agent is said to be of type 1 (respectively, type 2) if she has a unit-demand valuation with value 1 (respectively, β) for all items. Consider a market with m type-1 agents and $\lfloor \alpha \cdot m \rfloor$ type-2 agents. Consider the allocation that gives $\lfloor \alpha \cdot m \rfloor$ items to type-2 agents and $\lfloor (1 - \alpha) \cdot m \rfloor$ items to type-1 agents. This allocation gives a social welfare of $\lfloor \alpha \cdot m \rfloor \cdot \beta + \lfloor (1 - \alpha) \cdot m \rfloor \geq (\alpha \cdot m - 1) \cdot \beta + (1 - \alpha) \cdot m - 1 \geq \frac{\beta}{\beta-1}m - 4$. Given any pricing \mathbf{p} , let c be the number of prices in \mathbf{p} that are strictly greater than 1. We consider two cases:

- If $c > \frac{m}{\beta(\beta-1)} = (1 - \alpha)m$, then if the type-2 agents arrive first, followed by the type-1 agents, then there will not be any item left for type-1 agents. Therefore, the social welfare is at most

$$\alpha \cdot m \cdot \beta = m \cdot \frac{\beta + 1}{\beta} \leq 0.802 \cdot \left(\frac{\beta}{\beta - 1} m - 4 \right)$$

for large-enough m .

- If $c \leq \frac{m}{\beta(\beta-1)} = (1 - \alpha)m$, then if the type-1 agents arrive first, followed by the type-2 agents, then the best-case scenario is that the type-1 agents buy the $m - c$ cheap items while the type-2 agents buy the c expensive items. Therefore, the social welfare is at most

$$m - c + c \cdot \beta \leq m + \frac{m}{\beta} = m \cdot \frac{\beta + 1}{\beta} \leq 0.802 \cdot \left(\frac{\beta}{\beta - 1} m - 4 \right)$$

for large-enough m .

This concludes the argument. \square

The next result shows that if we know the order of the agents beforehand (while having no control over this order), then we can extract the full optimal welfare.

THEOREM 4.6. *For every market with submodular valuations with a known order of arrival, there exists a static pricing \mathbf{p} that guarantees the optimal social welfare.*

PROOF. Define b and ϵ as in Section 2, and number the agents from 1 to n according to their arrival order. Let x_1, \dots, x_n be the optimal allocation that allocates the items of marginal value b greedily (i.e., agent i gets an item of value b if and only if for all agents $j < i$, agent j also gets all items of marginal value b that she desires). Let k be the index of the last agent to get an item of marginal value b . Consider the pricing \mathbf{p} where $\sum_{i \leq k} x_i$ items are offered at price $b - \epsilon$ and $\sum_{i > k} x_i$ items are offered at price $b + \epsilon$. This pricing guarantees that agent i will buy exactly x_i items, and therefore guarantees the optimal social welfare as needed. \square

4.2 Uniform Pricing

In this section, we show that if we restrict ourselves to using uniform pricing with submodular valuations, we can still guarantee $1/2$ of the optimal welfare. Even though the result also follows from Theorem 8.2, which holds for XOS valuations in the Bayesian setting, we provide a simpler analysis for this particular case.

THEOREM 4.7. *For every market with submodular valuations, there exists a static uniform pricing p that guarantees at least $1/2$ of the optimal social welfare.*

PROOF. As in Theorem 4.2, setting a price of $b - \epsilon$ to all items guarantees a social welfare of at least mb . However, if we set a price of $b + \epsilon$ to all items, then exactly m' items will be sold—those items with marginal values greater than b —and their marginal values will be realized by the buyers. Thus, the social welfare is at least $OPT - (m - m')b$.

Therefore, using the best of the two prices guarantees at least $\max(mb, OPT - (m - m')b)$, which is at least half of OPT . \square

The last bound is tight, as shown in the following proposition.

PROPOSITION 4.8. *There exists a market with m items and two buyers with submodular valuations such that every uniform static pricing yields a social welfare of at most $\frac{m}{2m-1} (\approx \frac{1}{2})$ of the optimal social welfare.*

PROOF. Consider a market with m items and two buyers with the following valuations: Buyer 1 has a unit-demand valuation with a value of m for each item. Buyer 2 has an additive valuation with a value of 1 for each item. The optimal welfare is $OPT = 2m - 1$, obtained by allocating a single item to the first buyer, and $m - 1$ items to the second buyer. We show that no uniform pricing can guarantee more than a welfare of m . Let p be some uniform pricing. Distinguish between the following cases:

- (1) If $p \leq 1$, then if the additive buyer arrives first and buys all items, the social welfare is m .
- (2) If $p > 1$, then the additive buyer will never buy any item, and therefore the social welfare is at most m .

This concludes the argument. \square

The bound $1/2$ for uniform pricing is still tight even with identical buyers.

PROPOSITION 4.9. *There exists a market with buyers having identical submodular valuations such that every uniform static pricing yields a social welfare of at most $\frac{n+1}{2n} (\approx \frac{1}{2})$ of the optimal social welfare.*

PROOF. Consider a market with n^2 items and n buyers with the following valuation: $v(i) = n + i$ for $i \geq 1$. The optimal welfare is $OPT = 2n^2$, which can be achieved by giving each agent n items. We show that no uniform pricing can guarantee more than a welfare of $n^2 + n \approx OPT/2$. Let p be some uniform pricing. Distinguish between the following cases:

- (1) If $p \leq 1$, then the first buyer will buy all items, and therefore the social welfare is $n^2 + n$.
- (2) If $p > 1$, then every buyer will buy at most one item, and therefore the social welfare is at most $n \cdot (n + 1) = n^2 + n$.

This concludes the argument. \square

4.3 Dynamic Pricing

If we allow dynamic pricing, then the following result shows that we can extract the full optimal welfare.

THEOREM 4.10. *For every market with n agents with submodular valuations over m items, there exists a dynamic item pricing that guarantees the optimal social welfare.*

PROOF. We prove the theorem by showing a pricing that maintains the invariant that the current allocation can be completed to an optimal allocation in the induced market.

Define b and ϵ as in Section 2. Let $X \subseteq [n]$, $\ell \leq m$ be the remaining available agents and the number of available items in some iteration. Let k_i be the number of items worth strictly more than b to agent i , and let y_i be the number of items worth exactly b to agent i . The invariant that the allocation so far can be completed to an optimal allocation implies that $\sum_{i \in X} k_i \leq \ell$ and $\sum_{i \in X} (k_i + y_i) \geq \ell$. Consider the following pricing: If there exists an agent $i \in X$ such that $y_i > \ell - \sum_{j \in X} k_j$, then we set a price of $b - \epsilon$ to $\min_{i \in X: y_i > \ell - \sum_{j \in X} k_j} k_i + \ell - \sum_{j \in X} k_j$ items, and a price of $b + \epsilon$ to the remaining items. Otherwise, we set a price of $b - \epsilon$ to all items.

We claim that this pricing gives exactly OPT. Fix an iteration, and let i be the arriving buyer in that iteration. We show that:

- (1) Buyer i buys at least $\max(k_i, \ell - \sum_{j \in X, j \neq i} (k_j + y_j))$ items.
- (2) Buyer i buys at most $\min(\ell - \sum_{j \in X, j \neq i} k_j, k_i + y_i)$ items.
- (3) For every number of items $x \in [k_i + \max(0, \ell - \sum_{j \in X} k_j - \sum_{j \in X, j \neq i} y_j), \min(\ell - \sum_{j \in X, j \neq i} k_j, k_i + y_i)]$, it holds that $\text{OPT}(X \setminus \{i\}, \ell - x) + v_i(x) = \text{OPT}(X, \ell)$.

The combination of these properties ensures that the obtained allocation is optimal, since in each iteration, the value that agents i gets is equal to the loss of the social welfare for the remaining agents and items. We prove the three properties in turn.

- (1) First, if there exists $i' \in X$ such that $i' \neq i$ and $y_{i'} > \ell - \sum_{j \in X} k_j$, then $\ell - \sum_{j \in X, j \neq i} (k_j + y_j) \leq k_i + \ell - \sum_{j \in X} k_j - y_{i'} < k_i$, and therefore it is enough to show that agent i will buy at least k_i items. Agent i buys at least k_i items, since $\ell \geq \sum_{j \in X} k_j \geq k_i$ and therefore at least k_i items are available at price no more than $b + \epsilon$.

Second, if the only agent $i' \in X$ such that $y_{i'} > \ell - \sum_{j \in X} k_j$ is agent i , then we have $\min_{i' \in X: y_{i'} > \ell - \sum_{j \in X} k_j} k_{i'} = k_i$, and we get that the number of items that are offered at price $b - \epsilon$ is $k_i + \ell - \sum_{j \in X} k_j \leq k_i + y_i$. Therefore agent i will buy exactly $k_i + \ell - \sum_{j \in X} k_j \geq \max(k_i, \ell - \sum_{j \in X, j \neq i} (k_j + y_j))$ items, since $\ell - \sum_{j \in X} k_j \geq 0$.

Otherwise, there is no agent i' such that $y_{i'} > \ell - \sum_{j \in X} k_j$, and all items are priced at $b - \epsilon$. Therefore agent i buys $k_i + y_i \geq \ell - \sum_{j \in X, j \neq i} (k_j + y_j)$ items, since $\sum_{j \in X} (k_j + y_j) \geq \ell$.

- (2) If $y_i > \ell - \sum_{j \in X} k_j$, then $\min_{i' \in X: y_{i'} > \ell - \sum_{j \in X} k_j} k_{i'} \leq k_i$, and the number of items at price $b - \epsilon$ is at most $\ell - \sum_{j \in X, j \neq i} k_j$. Since $k_i \leq \ell - \sum_{j \in X, j \neq i} k_j$, agent i buys at most $\ell - \sum_{j \in X, j \neq i} k_j$ items. Otherwise, agent i buys at most $k_i + y_i \leq k_i + \ell - \sum_{j \in X} k_j = \ell - \sum_{j \in X, j \neq i} k_j$ items.

When all prices are at least $b - \epsilon$, agent i does not buy more than $k_i + y_i$ items.

- (3) We have

$$\text{OPT}(X, \ell) = \sum_{j \in X} v_j(k_j) + \left(\ell - \sum_{j \in X} k_j \right) \cdot b$$

$$\begin{aligned}
&= v_i(x) + \sum_{j \in X, j \neq i} v_j(k_j) + \left(\ell - x - \sum_{j \in X, j \neq i} k_j \right) \cdot b \\
&= v_i(x) + OPT(X \setminus \{i\}, \ell - x).
\end{aligned}$$

The second equation follows from the fact that $k_i \leq x \leq k_i + y_i$, implying that $v_i(x) = v_i(k_i) + (x - k_i)b$. For the last equation, we need to show that (a) $\ell - x - \sum_{j \in X, j \neq i} k_j \geq 0$ (i.e., there are enough items), and (b) $\ell - x - \sum_{j \in X, j \neq i} k_j \leq \sum_{j \in X, j \neq i} y_j$ (i.e., all items will be bought). Part (a) follows directly from (2) and part (b) follows directly from (1).

Therefore all three properties hold, and the proof is complete. \square

5 XOS VALUATIONS

In this section, we consider XOS valuations. We give upper bounds on the approximation ratio for both static and dynamic pricing.

THEOREM 5.1. *There exists a market of m items and two agents with XOS valuations for which no static pricing yields more than $1 - 1/e$ of the optimal social welfare.*

PROOF. Consider a market with m items for large-enough m . Let $k = \lfloor m/e \rfloor$. Assume the following agents' valuations:

$$\begin{aligned}
v_1(i) &= \begin{cases} k & i < k; \\ i & i \geq k; \end{cases} \\
v_2(i) &= \begin{cases} \sum_{0 \leq j < m-k} \frac{m-k-j}{m-1-j} & i > m-k; \\ \sum_{0 \leq j < i} \frac{m-k-j}{m-1-j} & i \leq m-k. \end{cases}
\end{aligned}$$

One can check that v_1 is XOS and v_2 is submodular. The allocation that maximizes the social welfare is the one that gives all items to agent 1, which generates a social welfare of m . We show that if agent 2 arrives first, then for every pricing \mathbf{p} , agent 1 will not buy more than one item. Assume agent 2 bought s items, so that all remaining items have a price of at least $\frac{m-k-s}{m-1-s}$ (otherwise, agent 2 would have bought more items). Agent 1's marginal utility from buying i items beyond a single item is

$$u_1(i+1) - u_1(1) \leq \begin{cases} -i \cdot \frac{m-k-s}{m-1-s} & i < k; \\ i - k - i \cdot \frac{m-k-s}{m-1-s} & m-s > i \geq k. \end{cases}$$

For the case where $i < k$, the marginal utility is clearly smaller than or equal to 0. For the case where $m-s > i \geq k$, the maximum marginal utility is obtained for $i = m-s-1$, where it is also smaller than or equal to 0. We conclude that agent 1 buys at most a single item if agent 2 arrives first. Therefore, the social welfare is bounded by $v_1(1) + v_2(m) = k + \sum_{0 \leq j < m-k} \frac{m-k-j}{m-1-j} = m - \sum_{0 \leq j < m-k} \frac{k-1}{m-1-j} = m - \sum_{k \leq j < m} \frac{k-1}{j} \approx m - k + 1 \approx (1 - 1/e) \cdot OPT$, where the last two equalities follow by the choice of $k = \lfloor m/e \rfloor$. \square

THEOREM 5.2. *There exists a market of three items and two agents with XOS valuations for which no dynamic pricing yields more than $5/6$ of the optimal social welfare.*

PROOF. Consider a market with two agents and three items, with the following valuations:

$$v_1(i) = \begin{cases} 4 & i \leq 2; \\ 6 & i = 3; \end{cases}$$

and the second agent has a unit-demand valuation with value 1. The allocation that maximizes the social welfare is the one that gives all items to agent 1, generating a social welfare of 6. Let \mathbf{p} be the prices in the first iteration. If there exists an item j with price $p_j \leq 1$, then if agent 2 arrives

first, she will buy item j , and the social welfare will be at most 5. If all items have prices greater than 1, then if agent 1 arrives first, then she will buy at most one item, and again the social welfare will be at most 5. \square

6 SUBADDITIVE VALUATIONS

In this section, we consider subadditive valuations. Recall that for submodular valuations (Section 4), we can set prices directly based on the values of b and ϵ . This method no longer works for subadditive valuations, since the marginal values of an agent do not necessarily form a non-increasing sequence. Nevertheless, we will show that it is still useful to consider the submodular closures of subadditive valuations that we introduced in Section 3 and set prices based on these closures.

Our main result of this section is the existence of a uniform price that guarantees at least 1/3 of the optimal welfare, which is the best possible using uniform pricing.

6.1 Upper Bounds

THEOREM 6.1. *For every market of m items with subadditive valuations, there exists a uniform static item pricing \mathbf{p} that guarantees at least 1/3 of the optimal social welfare.*

To prove Theorem 6.1, we first establish some technical lemmas. The first lemma provides a guarantee on the welfare obtained by an agent at a certain price in terms of the number of items left when the agent arrives.

LEMMA 6.2. *Let v be a subadditive valuation, and let \tilde{v} be the minimal submodular valuation such that $\tilde{v} \geq v$, as defined in Proposition 3.2. Let $b_k = \tilde{v}(k) - \tilde{v}(k-1)$ be the k th marginal value of \tilde{v} . When an agent with valuation v arrives, if there are k available items, all priced at $p = b_k/2$, then the welfare obtained from this agent is at least $kb_k/2$.*

PROOF. For notational simplicity, we write $b = b_k$. Let $\ell = \max_i \{\tilde{v}(i) - \tilde{v}(i-1) \geq b\}$; notice that $\ell \geq k$ by definition. Let α be the integer such that $\frac{\ell}{\alpha} \leq k < \frac{\ell}{\alpha-1}$. We can assume without loss of generality that $\alpha \geq 2$, since for $\alpha = 1$ the agent will buy $k (= \ell)$ items and $v(k) = \tilde{v}(k) \geq kb/2$. Let $c = \max_i \{\tilde{v}(i) - \tilde{v}(i-1) > b\}$; notice that $c < k$ by submodularity of \tilde{v} .

Let $u(x) = v(x) - x \cdot b/2$ be the utility that an agent with valuation v derives from obtaining x items. Let $\underline{a} = \lfloor \frac{\ell}{\alpha} \rfloor$, $\bar{a} = \lceil \frac{\ell}{\alpha} \rceil$, and $\beta \equiv \ell \pmod{\alpha}$ with $0 \leq \beta < \alpha$. By simple algebra it holds that

$$(\alpha - \beta) \cdot \underline{a} + \beta \cdot \bar{a} = \ell. \quad (4)$$

We have:

$$\begin{aligned} (\alpha - \beta) \cdot u(\underline{a}) + \beta \cdot u(\bar{a}) &= (\alpha - \beta) \cdot \left(v(\underline{a}) - \underline{a} \cdot \frac{b}{2} \right) + \beta \cdot \left(v(\bar{a}) - \bar{a} \cdot \frac{b}{2} \right) \\ &= (\alpha - \beta) \cdot v(\underline{a}) + \beta \cdot v(\bar{a}) - \beta \cdot \bar{a} \cdot \frac{b}{2} - (\alpha - \beta) \cdot \underline{a} \cdot \frac{b}{2} \geq v(\ell) - \frac{\ell b}{2}, \end{aligned}$$

where the last inequality follows from the subadditivity of v and Equation (4). Since the first expression is a sum of α terms, there exists a number of items $x' \in \{\underline{a}, \bar{a}\}$ such that $u(x') \geq \frac{v(\ell) - \frac{\ell b}{2}}{\alpha}$. Now, let x^* be the number of items that the agent actually buys. We know that $x^* \geq c$, since the first c items yield an average marginal value of more than b . Moreover, since the agent buys x^* items when she has a choice of buying x' , we have $u(x^*) \geq u(x') \geq \frac{v(\ell) - \frac{\ell b}{2}}{\alpha}$. We now consider two cases:

- (1) $v(c) \geq kb/2$. In this case, since we have that $x^* \geq c$ it follows immediately that $v(x^*) \geq v(c) \geq kb/2$ as needed.

- (2) $v(c) < kb/2$. Since all marginal values beyond the c th item are no greater than b , we have $v(c+x) \leq v(c) + bx$ for any $x \geq 0$. This implies that $u(c+x) \leq u(c) + xb/2$. In particular, since $x^* \geq c$ and x^* yields the optimal utility to the agent, we have $u(c) \leq u(x^*) \leq u(c) + (x^* - c)b/2$. This implies that $x^* \geq c + \left\lceil \frac{u(x^*) - u(c)}{b/2} \right\rceil$, and it follows that:

$$\begin{aligned}
v(x^*) &= v(c) + (x^* - c) \cdot \frac{b}{2} + u(x^*) - u(c) \\
&\geq v(c) + \left\lceil \frac{u(x^*) - u(c)}{b/2} \right\rceil \cdot \frac{b}{2} + u(x^*) - u(c) \\
&\geq v(c) + u(x^*) - u(c) + u(x^*) - u(c) \\
&= v(c) + 2u(x^*) - 2u(c) \\
&\geq v(c) + 2 \cdot \frac{v(\ell) - \frac{\ell b}{2}}{\alpha} - 2 \left(v(c) - \frac{cb}{2} \right) \\
&= \frac{2v(\ell) - \ell b - \alpha v(c) + \alpha cb}{\alpha} \\
&= \frac{2(v(c) + (\ell - c)b) - \ell b - \alpha v(c) + \alpha cb}{\alpha} \\
&= \frac{v(c)(2 - \alpha) + cb(\alpha - 2) + \ell b}{\alpha} \\
&\geq \frac{\frac{kb}{2}(2 - \alpha) + \ell b}{\alpha} \\
&\geq \frac{\frac{\ell b}{2(\alpha - 1)}(2 - \alpha) + \ell b}{\alpha} \\
&= \frac{\ell b(2 - \alpha) + 2\ell b(\alpha - 1)}{2\alpha(\alpha - 1)} \\
&= \frac{\ell b}{2(\alpha - 1)} \\
&\geq \frac{kb}{2},
\end{aligned}$$

as needed.

This completes the argument. \square

Since we want to set the same price for all items, we need a way to compare the welfare obtained from setting different prices to apply Lemma 6.2. The next lemma states that the number of items that an agent buys can only decrease as the price goes up.

LEMMA 6.3. *Let v be any non-decreasing valuation, and let $p_1 < p_2$ be two prices (assigned to all items). Let $x_1 = \operatorname{argmax}(v(x) - p_1 \cdot x)$, and $x_2 = \operatorname{argmax}(v(x) - p_2 \cdot x)$. Then $x_1 \geq x_2$.*

PROOF. We know that $v(x_1) - x_1 \cdot p_1 \geq v(x_2) - x_2 \cdot p_1$ and $v(x_2) - x_2 \cdot p_2 \geq v(x_1) - x_1 \cdot p_2$. Using these inequalities gives us:

$$\begin{aligned}
\frac{v(x_1) - v(x_2)}{p_1} &\geq x_1 - x_2 \geq \frac{v(x_1) - v(x_2)}{p_2} \\
\Rightarrow (p_2 - p_1)(v(x_1) - v(x_2)) &\geq 0
\end{aligned}$$

$$\begin{aligned} &\Rightarrow v(x_1) - v(x_2) \geq 0 \\ &\Rightarrow v(x_1) \geq v(x_2). \end{aligned}$$

If it were the case that $x_1 < x_2$, then since v is non-decreasing, we must have $v(x_1) = v(x_2)$. Our initial inequalities then imply that $x_1 = x_2$, a contradiction. So $x_1 \geq x_2$, as claimed. \square

By combining Lemmas 6.2 and 6.3, we can derive a guarantee on the social welfare obtained by setting a uniform price.

LEMMA 6.4. *Let v_1, \dots, v_n be the subadditive valuations of the agents, and let $\tilde{v}_1, \dots, \tilde{v}_n$ be the respective minimal submodular valuations such that $\tilde{v}_i \geq v_i$, as defined in Proposition 3.2. Define b as in Section 2 with respect to the \tilde{v}_i 's. When offering the agents the price $b/2$, the social welfare obtained will be at least $mb/2$.*

PROOF. If all items were sold, then the revenue will be $mb/2$ and therefore the social welfare will be at least $mb/2$ as needed. Otherwise, let $\ell_i = \max\{x \mid (\tilde{v}_i(x) - \tilde{v}_i(x-1)) \geq b\}$. We know that there exists an agent i such that she bought less than ℓ_i items (else all items would have been sold). Suppose that upon the arrival of agent i there are k items remaining for sale. If $k \geq \ell_i$, then agent i will buy at least ℓ_i items, contradicting our assumption, so $k < \ell_i$.

Let b' be the k th marginal value of \tilde{v}_i ; in particular, $b' \geq b$. By Lemma 6.2, if we set a price of $b'/2$, then the welfare obtained from agent i is at least $kb'/2$. Now, by Lemma 6.3, we can only sell more items to agent i by setting a lower price $b/2$. This implies that the welfare obtained by this agent is at least $kb/2$. In addition, it is known that $m - k$ items were sold prior to the arrival of agent i , contributing at least $(m - k)b/2$ to the social welfare. We conclude that the total social welfare is at least $(m - k)b/2 + kb/2 = mb/2$. \square

With Lemma 6.4 in hand, we can now derive the 1/3 approximation on the social welfare with uniform pricing.

PROOF OF THEOREM 6.1. Let v_1, \dots, v_n be the subadditive valuations of the agents, and let $\tilde{v}_1, \dots, \tilde{v}_n$ be the respective minimal submodular valuations such that $\tilde{v}_i \geq v_i$, as defined in Proposition 3.2. Define V, b, m' , and ϵ with respect to the \tilde{v}_i 's as in Section 2. We show that one of the following pricing schemes guarantees at least 1/3 of the optimal social welfare:

- (P1) a uniform price of $b + \epsilon$;
- (P2) a uniform price of $b/2$.

Let (x_1, \dots, x_n) be the optimal partition of items to agents (where x_i is the number of items allocated to agent i), and let $OPT = v_1(x_1) + \dots + v_n(x_n)$ be the optimal social welfare. It is easy to see that $OPT \leq \sum_{b < b' \in V} b' + (m - m')b$ (since the RHS is the optimal allocation under valuations \tilde{v}_i 's, which is at least as high as the optimal welfare of v_i 's).

Pricing (P1) guarantees exactly a welfare of $\sum_{b < b' \in V} b'$. Pricing (P2) by Lemma 6.4 guarantees a welfare of at least $mb/2$. The best of the two guarantees yields at least 1/3 of the optimal social welfare. \square

If there are two identical agents, then the bound in Theorem 6.1 can be improved to 2/3.

THEOREM 6.5. *For every market of m items and two identical agents with subadditive valuations, there exists a uniform static item pricing \mathbf{p} that guarantees at least 2/3 of the optimal social welfare.*

PROOF. Let OPT be the optimal social welfare, and v the valuation of the two agents. Set a uniform price of $OPT/3m$. In the optimal allocation, the total utility of the two agents, if they were to buy the items at this price, is $OPT - m \cdot OPT/3m = 2OPT/3$, so at least one agent gets utility

at least $OPT/3$. This means that in the allocation that results from setting the price $OPT/3m$, the first agent must also get utility at least $OPT/3$, since all items are still available. If the first agent buys at most half of the items, then so does the second agent, and the total utility (and hence the welfare) is at least $2OPT/3$ as we want. So suppose that the first agent buys $m' > m/2$ items. The first agent contributes $m' \cdot OPT/3m$ to the revenue and has utility at least $OPT/3$, so since welfare is the sum of revenue and utility, we have $v(m') \geq \frac{OPT}{3} \cdot \left(1 + \frac{m'}{m}\right)$.

It suffices to show that the second agent contributes at least $2OPT/3 - v(m')$ to the welfare. We consider two cases:

- (1) $v(m - m') \geq \frac{2OPT}{3} - v(m') + \frac{OPT}{3m} \cdot (m - m')$. In this case, the second agent can guarantee utility at least $2OPT/3 - v(m')$ by buying all of the remaining $m - m'$ items, so the agent's actual utility (and hence contribution to the welfare) is at least $2OPT/3 - v(m')$.
- (2) $v(m - m') < \frac{2OPT}{3} - v(m') + \frac{OPT}{3m} \cdot (m - m')$. Let $k := \lfloor \frac{m'}{m-m'} \rfloor$, and let $x := m' - k(m - m') = (k + 1)m' - km \geq 0$. By considering the possibility that the second agent buys x items, we get the desired conclusion if $v(x) \geq \frac{2OPT}{3} - v(m') + \frac{OPT}{3m} \cdot x$. Hence, we may assume that $v(x) < \frac{2OPT}{3} - v(m') + \frac{OPT}{3m} \cdot x$. We have

$$\begin{aligned} v(m') &= v(k(m - m') + x) \\ &\leq k \cdot v(m - m') + v(x) \\ &< (k + 1) \left(\frac{2OPT}{3} - v(m') \right) + \frac{OPT}{3m} \cdot m', \end{aligned}$$

where the first inequality follows from subadditivity of v . Rearranging, we find that $(k + 2) \cdot v(m') < \frac{OPT}{3} \cdot \left(2k + 2 + \frac{m'}{m}\right)$. Combining this with the previous estimate $v(m') \geq \frac{OPT}{3} \cdot \left(1 + \frac{m'}{m}\right)$, we have $(k + 2) \left(1 + \frac{m'}{m}\right) < 2k + 2 + \frac{m'}{m}$, which rearranges to $x < 0$, a contradiction. Hence the second case cannot occur.

This completes the proof. \square

6.2 Lower Bounds

The next propositions show that the bound in Theorem 6.1 cannot be improved to more than $1/2$, and in the case of using only uniform pricing, cannot be improved to more than $1/3$. This means that the bound in Theorem 6.1 is tight for uniform pricing.

PROPOSITION 6.6. *There is a market with subadditive valuations with m items and two agents such that no static pricing \mathbf{p} guarantees more than $1/2$ of the optimal social welfare.*

PROOF. Consider a market where the first agent has a valuation $v_1(i) = 1$ for $i \leq m - 1$ and $v_1(m) = 2$, and the second agent has a unit-demand valuation with value $\frac{1}{m-1}$. The optimal allocation is giving all the items to the first agent (therefore the optimal social welfare is 2). Let \mathbf{p} be any pricing. If there exists i such that $p_i \leq \frac{1}{m-1}$, then if agent 2 arrives first, she will buy an item and the social welfare will be bounded by $1 + \frac{1}{m-1}$. Otherwise, we have $p_i > \frac{1}{m-1}$ for all i . In this case, agent 1 has no incentive to buy more than one item, and again the social welfare will be bounded by $1 + \frac{1}{m-1}$. Taking m large, this yields the ratio of $1/2$ as claimed. \square

PROPOSITION 6.7. *There is a market with subadditive valuations with m items and three agents such that no uniform static pricing \mathbf{p} guarantees more than $1/3$ of the optimal social welfare.*

PROOF. Consider a market where the first agent has a valuation $v_1(i) = m - 1 + i$ for $i \leq m - 1$ and $v_1(m) = 3m - 2$, the second agent has a unit-demand valuation with value 2, and the third agent has an additive valuation $v_3(i) = i$. The optimal allocation is giving all items to the first

agent, which yields the optimal social welfare of $3m - 2$. Let p be any uniform pricing. If $p \leq 1$, then if agent 3 arrives first, she will buy all the items and the social welfare will be bounded by m . If $1 < p \leq 2$, then if agent 2 arrives first, then she will buy an item and then only the first agent will buy an additional item, and the social welfare will be bounded by $m + 2$.

Otherwise, we have that $p > 2$. In this case, agent 1 has no incentive to buy more than one item, the other agents will not buy any item, and again the social welfare will be bounded by m . Taking m large, this yields a ratio of $1/3$, as claimed. \square

In the case of two identical agents, the approximation cannot be improved to more than $3/4$. In this special case, we can guarantee at least half of the social welfare by applying Theorem 7.3.

PROPOSITION 6.8. *There exists a market of m items and two identical agents with a subadditive valuation such that no static pricing guarantees more than $3/4$ of the optimal social welfare.*

PROOF. Assume that m is even for simplicity. Consider the valuation

$$v(i) = \begin{cases} m/4 & i < m/2 \\ m/2 - 1 - \epsilon & i = m/2, \\ m/2 & i > m/2 \end{cases}$$

and consider a market where there are two agents with valuation v . The optimal allocation is to give $m/2$ items to each agent, which generates a social welfare of $m - 2 - 2\epsilon$. To set prices so that each agent will want to buy exactly $m/2$ items, we need that the $m/2$ more expensive items cost less than $m/2 - 1 - \epsilon$, and therefore there exists one among them that costs less than 1. When the first agent arrives, she will buy either exactly $m/2 + 1$ items or less than $m/2$ items, since the cheapest $m/2 + 1$ items cost less than 1 each. Therefore the social welfare obtained by any static pricing in this market cannot be more than $3m/4 \approx 3OPT/4$, as claimed. \square

If we use uniform pricing, then we cannot guarantee more than $2/3$ of the welfare for two identical agents. This means that the bound in Theorem 6.5 is tight.

PROPOSITION 6.9. *There is a market with subadditive valuations with m items and two identical agents such that no uniform static pricing \mathbf{p} guarantees more than $2/3$ of the optimal social welfare.*

PROOF. Consider a market where both agents have the valuation $v(i) = m + i$. An optimal allocation is to give one item to the first agent and the remaining items to the second agent for a social welfare of $3m$. Let p be any uniform pricing. If $p \leq 1$, then the first agent buys all the items and the social welfare is $2m$. Otherwise, we have $p > 1$, and both agents only buy one item for a social welfare of $2m + 2$. Taking m large, this yields a ratio of $2/3$ as claimed. \square

7 GENERAL VALUATIONS

In this section, we consider general valuations. We show that even in the case of two agents, the welfare achieved by posted prices can be a linear factor away from the optimal social welfare. However, we establish a constant guarantee of $1/2$ if we are allowed to choose the order in which the agents arrive. Perhaps surprisingly, even for general valuations it will still be helpful to consider the submodular closures that we defined in Section 3.

While the analysis assumes monotonicity, all results hold even for non-monotone valuations: simply do all calculations based on the monotone closures of the valuations, and the results hold for the valuations themselves.

7.1 Worst-case Ordering

We first show that for general valuations, we cannot guarantee more than $1/m$ of the optimal welfare even if we know the order of arrival.

PROPOSITION 7.1. *There is a market with valuations over m items and two agents such that no static pricing \mathbf{p} guarantees more than $1/m$ of the optimal social welfare even for a known order of arrival.*

PROOF. Consider a market where the first agent has a unit-demand valuation with value 1, and the second agent is single-minded for the set of all items with value m . The optimal allocation is giving all the items to the second agent (therefore the optimal social welfare is m). Let \mathbf{p} be any pricing. If there exists i such that $p_i < 1$, then the first agent will buy an item and the social welfare will be bounded by 1. Otherwise, $p_i \geq 1$ for all i , and therefore the price of the bundle of all items is at least m . This means agent 2 has no incentive to buy the bundle and the social welfare obtained by these prices is 0. \square

The example in Proposition 7.1 is as bad as it gets: We can guarantee $1/m$ of the optimal welfare with uniform pricing.

THEOREM 7.2. *For every market of m identical items, there exists a uniform static item pricing \mathbf{p} that guarantees at least $1/m$ of the optimal social welfare.*

PROOF. Consider an allocation that yields the optimal welfare, and let β be the highest average welfare per item among all agents in this allocation. Clearly, $OPT \leq m\beta$. Set the uniform price to be $\beta - \epsilon$ for small-enough ϵ . Then at least one item will be sold, since the agent with the highest average welfare per item in the optimal allocation can achieve positive utility. It follows that the welfare obtained by this pricing is at least β . \square

7.2 Best-case Ordering

Next, we show that if we can choose the order of arrival, then we can guarantee at least half of the optimal welfare. We remark that when agents are identical, the order of arrival does not matter, and therefore our result holds for the setting with identical agents as well.

THEOREM 7.3. *For every market of m identical items, there exists a uniform static item pricing \mathbf{p} along with an order of arrival that guarantees at least $1/2$ of the optimal social welfare.*

PROOF. Let v_1, \dots, v_n be the valuations of the agents, and let $\tilde{v}_1, \dots, \tilde{v}_n$ be the respective minimal submodular valuations such that $\tilde{v}_i \geq v_i$, as defined in Proposition 3.2. Define V, b , and m' as in Section 2 with respect to the \tilde{v}_i 's. In addition, one can find a sufficiently small ϵ so that at the price $b - \epsilon$, whenever an agent buys an additional set of items that belong to the same "marginal value segment" of \tilde{v}_i , the average marginal value of items in the set is at least b . Let (x_1, \dots, x_n) be the optimal allocation of items to agents where x_i is the number of items allocated to agent i , and let $OPT = v_1(x_1) + \dots + v_n(x_n)$ be the optimal social welfare.

We show how to construct a pricing scheme that guarantees at least $OPT/2$:

- (P1) a uniform price of $b + \epsilon$ with any arrival order (in this case the order does not influence the social welfare);
- (P2) a uniform price of $b - \epsilon$ with some arrival order σ to be chosen.

It is easy to see that $OPT \leq \sum_{b < b' \in V} b' + (m - m')b$, since the RHS is the optimal allocation under valuations \tilde{v}_i 's, which is at least as high as the optimal welfare under v_i 's.

Pricing (P1) guarantees a welfare of exactly $\sum_{b < b' \in V} b'$. We next show that we can generate orders σ, σ' so that the social welfare of the best of pricing (P1) and pricing (P2) with order σ or σ' will guarantee at least $OPT/2$. We start with an arbitrary order σ for (P2) and consider three cases.

- (1) All items are sold. The social welfare will be at least mb , and the best of pricing (P1) and (P2) with order σ will generate a social welfare of at least $OPT/2$.
- (2) Not all items are sold, but the number of items bought by every agent i is at least the number of items with marginal contribution greater than b (with respect to \tilde{v}_i). Suppose that $m' \leq k < m$ items are sold in total (and therefore the social welfare of pricing (P2) and order σ is exactly $\sum_{b < b' \in V} b' + (k - m')b$). In this case, we have that there exists an agent i such that $\tilde{v}_i(m - k + 1) - \tilde{v}_i(m - k) \geq b$; otherwise, the agents would get all the items that contribute to them at least b until at least $k + 1$ items are sold. We set σ' to be an order where agent i is the first agent to arrive. In order σ' and pricing (P2), we have that the first agent (i.e., agent i) will buy at least $m - k + 1$ items (since the marginal value of the first $m - k + 1$ items is at least b and the price is less than b , and all items are available in this stage). The sum of the social welfare of the two schemes (with order σ or σ') is at least $\sum_{b < b' \in V} b' + (k - m')b + (m - k + 1)b \geq OPT$, and therefore the best of the two will generate at least $OPT/2$ as needed.
- (3) Not all items are sold, and the number of items bought by some agent i is less than the number of items with marginal contribution greater than b (with respect to \tilde{v}_i). Suppose that k items are sold in total. In this case, we have that $k > m - m'$; otherwise, when an agent with marginal value greater than b is offered, there will be enough items available for her to buy all items with marginal value greater than b (since there are at most m' such items and at least m' items are available for her to buy). Therefore the social welfare of this pricing scheme is at least $kb > (m - m')b$. Hence the best of this pricing scheme and pricing scheme (P1) will generate a social welfare of at least $OPT/2$.

This concludes the argument. □

The bound $1/2$ in Theorem 7.3 is also tight.

PROPOSITION 7.4. *There exists a market of m items for which no static pricing and order of arrival yields more than $1/2$ of the optimal social welfare.*

PROOF. Assume that m is even for simplicity. Consider the valuation

$$v(i) = \begin{cases} 0 & i < m/2 \\ m/2 - \epsilon & i = m/2. \\ m/2 + 1 & i > m/2 \end{cases}$$

And consider a market where there are two agents with valuation v . Since the agents have the same valuation, the order of arrival does not influence the social welfare. The optimal allocation is to give $m/2$ items to each agent. To set prices so that each agent will want to buy at least $m/2$ items, we need that the $m/2$ more expensive items together cost less than $m/2 - \epsilon$, and therefore there exists one among them that costs less than 1. When the first agent arrives, she will buy exactly $m/2 + 1$ items, since the cheapest $m/2 + 1$ cost less than 1 each. Therefore the social welfare obtained by any static pricing in this market cannot exceed $m/2 + 1 \approx OPT/2$, as claimed. □

8 BAYESIAN SETTING

In this section, we consider the Bayesian setting, where the valuation of each agent is drawn independently from a distribution that can be different for different agents. Given the generality of this setting, the methods that we develop for the full-information setting no longer apply; instead, we will rely mainly on the techniques from previous work by Feldman et al. [2015].

8.1 XOS Valuations

Feldman et al. [2015] showed that if agents' valuations are drawn independently from a distribution over XOS valuations, then there exist prices that yield expected welfare at least half of the expected optimal welfare. These posted prices are non-uniform, even when the result is specialized to the case of identical items. We first restate Feldman et al.'s result and then show that if the items are identical, then the same bound can be obtained using uniform prices.

THEOREM 8.1 [FELDMAN ET AL. 2015]. *Let $\mathcal{F} = F_1 \times \dots \times F_n$ be a product distribution over XOS valuations (not necessarily over identical items). For every $\mathbf{v} = (v_1, \dots, v_n) \in \mathcal{F}$, let $X^*(\mathbf{v}) = (X_1^*(\mathbf{v}), \dots, X_n^*(\mathbf{v}))$ be any allocation that maximizes the social welfare. Let $\mathbf{a} = (a_1, \dots, a_n)$ be additive valuations such that $v_i(S) \geq a_i(S)$ for any subset S of items, and $v_i(X_i^*) = a_i(X_i^*)$. When the items are offered at prices $p_j = E_{\mathbf{v} \in \mathcal{F}} [a_i(j)/2]$, where $j \in X_i^*(\mathbf{v})$, the expected social welfare is at least $OPT/2$.*

THEOREM 8.2. *Let $\mathcal{F} = F_1 \times \dots \times F_n$ be a product distribution over XOS valuations (over identical items). Let OPT be the expected optimal social welfare. When all items are offered at the uniform price $OPT/2m$, the expected social welfare is at least $OPT/2$.*

PROOF. Let G be the uniform distribution over m different classes of the zero valuation. We notate the m different classes as 0_σ for $\sigma \in [m]$. Let $\hat{\mathcal{F}}$ be the distribution $\mathcal{F} \times G$. In other words, in addition to the distribution \mathcal{F} over the valuations of the n agents, we also add another agent whose valuation is always zero. Clearly, the distributions $\hat{\mathcal{F}}$ and \mathcal{F} have the same expected optimal social welfare, and for any prices $\mathbf{p} > 0$, the expected social welfare when considering $\hat{\mathcal{F}}$ and \mathcal{F} is the same, since the added agent never buys an item.

Let $\mathbf{v} = (v_1, \dots, v_n) \in \mathcal{F}$, and let $X^*(\mathbf{v}) = (X_1^*(\mathbf{v}), \dots, X_n^*(\mathbf{v}))$ be any optimal allocation with respect to \mathbf{v} . We define an optimal allocation for every $\hat{\mathbf{v}} = (v_1, v_2, \dots, v_n, 0_\sigma) \in \hat{\mathcal{F}}$ to be $\hat{X}_i^*(\hat{\mathbf{v}}) = \{j \in [m] \mid j + \sigma \in X_i^*(\mathbf{v})\}$, where $j + \sigma$ is considered modulo m . This is an optimal allocation, since the valuations are over identical items and the allocation gives the same number of items to each agent as the original optimal allocation. Let $\mathbf{a} = (a_1, \dots, a_n)$ be the additive valuations corresponding to $X^*(\mathbf{v})$ that maximize social welfare, as defined in Theorem 8.1. Let $\hat{\mathbf{a}} = (\hat{a}_1, \dots, \hat{a}_n)$ be the additive valuations where $\hat{a}_i(\{j\}) = a_i(\{j + \sigma\})$. Due to the symmetry over the items, these valuations also maximize social welfare.

We now show that with the new distribution $\hat{\mathcal{F}}$, the prices are all identical and equal to $OPT/2m$. The prices calculated according to Theorem 8.1 are

$$\begin{aligned}
 p_j &= E_{\hat{\mathbf{v}} \sim \hat{\mathcal{F}}} [\hat{a}_i(\{j\})/2 \text{ where } j \in \hat{X}_i^*(\hat{\mathbf{v}})] \\
 &= E_{\mathbf{v} \sim \mathcal{F}, 0_\sigma \sim G} [a_i(\{j + \sigma\})/2 \text{ where } j + \sigma \in X_i^*(\mathbf{v})] \\
 &= \frac{1}{2m} \cdot E_{\mathbf{v} \sim \mathcal{F}} \left[\sum_{j'} a_i(\{j'\}) \text{ where } j' \in X_i^*(\mathbf{v}) \right] \\
 &= \frac{1}{2m} \cdot OPT(\mathbf{v}) \\
 &= \frac{OPT}{2m}.
 \end{aligned}$$

This concludes the argument. □

8.2 Subadditive and General Valuations

We now define a notion that describes how close an arbitrary valuation is to an XOS valuation and derive approximation results in terms of this closeness quantity. The proof of Theorem 8.3 follows the analysis presented by Feldman et al. [2015].

Definition 8.1. We say that a valuation v (not necessarily over identical items) is c -close to XOS if there exists an XOS valuation \tilde{v} such that for every set of item S , it holds that $v(S)/c \leq \tilde{v}(S) \leq v(S)$.

THEOREM 8.3. *For any product distribution \mathcal{F} over valuations (not necessarily over identical items) that are c -close to XOS, there exist anonymous prices \mathbf{p} that guarantee an expected social welfare of at least $1/(2c)$ of the optimal expected welfare.*

PROOF. For every $\mathbf{v} = (v_1, \dots, v_n)$, let $X^*(\mathbf{v}) = (X_1^*, \dots, X_n^*)$ be an allocation that maximizes the social welfare with respect to the valuations \mathbf{v} . Let $\tilde{\mathbf{v}} = (\tilde{v}_1, \dots, \tilde{v}_n)$ be the vector of the corresponding XOS valuations; i.e., $v_i/c \leq \tilde{v}_i \leq v_i$ for every i . Let $\mathbf{a} = (a_1, \dots, a_n)$ be additive valuations such that $\tilde{v}_i(S) \geq a_i(S)$ for any subset S of items, and $\tilde{v}_i(X_i^*) = a_i(X_i^*)$.

For every item j , we define w_j to be the social welfare contribution of item j under valuations \mathbf{a} , and we set its price to be $p_j = E_{\mathbf{v} \sim \mathcal{F}}[w_j]/2$. (Recall that every vector of valuations \mathbf{v} defines an optimal allocation $X^*(\mathbf{v})$ and XOS valuations $\tilde{\mathbf{v}}$, and the allocation and the XOS valuations in turn define additive valuations \mathbf{a}). The expected maximal social welfare is $E[\text{SW}] = E_{\mathbf{v} \sim \mathcal{F}} \left[\sum v_i(X_i^*) \right]$.

Fix i and (v_i, \mathbf{v}_{-i}) . Let $\text{SOLD}_i(\mathbf{v})$ denote the items that have been sold before the arrival of agent i . Consider $\mathbf{v}'_{-i} \sim \mathcal{F}_{-i}$, which is independent of \mathbf{v} , and let $\tilde{\mathbf{v}}'_{-i}$ be the XOS valuations that approximate \mathbf{v}'_{-i} . Let $X^*(v_i, \mathbf{v}'_{-i})$ be an optimal allocation under the valuations (v_i, \mathbf{v}'_{-i}) ; in particular, $X_i^*(v_i, \mathbf{v}'_{-i})$ is the bundle assigned to agent i in this allocation. Define \mathbf{a}' analogously to \mathbf{a} with respect to the allocation $X^*(v_i, \mathbf{v}'_{-i})$ and the valuations $\tilde{\mathbf{v}}'_{-i}$. Let $S_i(v_i, \mathbf{v}_{-i}, \mathbf{v}'_{-i}) = X_i^*(v_i, \mathbf{v}'_{-i}) \setminus \text{SOLD}_i(\mathbf{v})$ be the subset of items in $X_i^*(v_i, \mathbf{v}'_{-i})$ that are available when agent i arrives; therefore her utility is at least the utility from purchasing this set of items. Thus her utility satisfies

$$u_i(\mathbf{v}) \geq E_{\mathbf{v}'_{-i}} \left[\sum_{j \in S_i(v_i, \mathbf{v}_{-i}, \mathbf{v}'_{-i})} (a'(\{j\}) - p_j)^+ \right],$$

where $f^+ = \max(f, 0)$. Here we use the assumption that $a'_i(S) \leq \tilde{v}_i(S) \leq v_i(S)$ for any subset S of items. Adding these inequalities for all i and taking expectation over \mathbf{v} gives us:

$$\begin{aligned} E_{\mathbf{v} \sim \mathcal{F}} \left[\sum_{i=1}^n u_i(\mathbf{v}) \right] &\geq \sum_{j \in M} \sum_{i=1}^n E_{v_i, \mathbf{v}_{-i}, \mathbf{v}'_{-i}} [1_{j \in X_i^*(v_i, \mathbf{v}'_{-i})} \cdot (a'(\{j\}) - p_j)^+ \cdot 1_{j \notin \text{SOLD}_i(\mathbf{v})}] \\ &= \sum_{j \in M} \sum_{i=1}^n E_{v_i, \mathbf{v}_{-i}, \mathbf{v}'_{-i}} [1_{j \in X_i^*(v_i, \mathbf{v}'_{-i})} \cdot (a'(\{j\}) - p_j)^+ \cdot 1_{j \notin \text{SOLD}_i(\mathbf{v}_{-i})}] \\ &= \sum_{j \in M} \sum_{i=1}^n E_{\mathbf{v}_{-i}} [1_{j \notin \text{SOLD}_i(\mathbf{v}_{-i})}] \cdot E_{v_i, \mathbf{v}'_{-i}} [1_{j \in X_i^*(v_i, \mathbf{v}'_{-i})} \cdot (a'(\{j\}) - p_j)^+] \\ &\geq \sum_{j \in M} \sum_{i=1}^n \Pr_{\mathbf{v}} [j \notin \text{SOLD}(\mathbf{v})] \cdot E_{v_i, \mathbf{v}'_{-i}} [1_{j \in X_i^*(v_i, \mathbf{v}'_{-i})} \cdot (a'(\{j\}) - p_j)^+] \\ &= \sum_{j \in M} \left(\Pr_{\mathbf{v}} [j \notin \text{SOLD}(\mathbf{v})] \cdot \sum_{i=1}^n E_{\mathbf{v}} [1_{j \in X_i^*(\mathbf{v})} \cdot (a(\{j\}) - p_j)^+] \right) \end{aligned}$$

$$\begin{aligned}
&\geq \sum_{j \in M} \left(Pr_{\mathbf{v}}[j \notin SOLD(\mathbf{v})] \cdot \sum_{i=1}^n E_{\mathbf{v}}[1_{j \in X_i^*(\mathbf{v})} \cdot (a(\{j\}) - p_j)] \right) \\
&= \sum_{j \in M} Pr_{\mathbf{v}}[j \notin SOLD(\mathbf{v})] \cdot p_j.
\end{aligned}$$

In the second line, we replace $SOLD_i(\mathbf{v})$ by $SOLD_i(\mathbf{v}_{-i})$, since the set of items sold before the arrival of agent i does not depend on v_i . In the fourth line, we decrease each probability $Pr_{\mathbf{v}_{-i}}[j \notin SOLD(\mathbf{v}_{-i})]$ to $Pr_{\mathbf{v}}[j \notin SOLD(\mathbf{v})]$; the inequality holds, since all the remaining terms are non-negative. In the fifth line, we replace \mathbf{v}'_{-i} by \mathbf{v}_{-i} , since these valuations are drawn from the same distribution and we are taking their expectation. The last line follows from the definition of p_j .

The expected revenue of this mechanism is $\sum_j Pr_{\mathbf{v}}[j \in SOLD(\mathbf{v})] \cdot p_j$. Therefore the expected social welfare guaranteed by these prices is at least $\sum_{j \in M} Pr_{\mathbf{v}}[j \in SOLD(\mathbf{v})] \cdot p_j + \sum_{j \in M} Pr_{\mathbf{v}}[j \notin SOLD(\mathbf{v})] \cdot p_j = \sum_{j \in M} p_j$. Finally,

$$\sum_{j \in M} p_j = \frac{1}{2} \cdot E_{\mathbf{v}} \left[\sum_{j \in M} w_j \right] = \frac{1}{2} \cdot E_{\mathbf{v}} \left[\sum_{i=1}^n \tilde{v}_i(X_i^*) \right] \geq \frac{1}{2c} \cdot E_{\mathbf{v}} \left[\sum v_i(X_i^*) \right],$$

and therefore we have a $2c$ approximation, as claimed. \square

Since any subadditive valuation is 2 -close to XOS according to Proposition 3.4, Theorem 8.3 implies that we can obtain at least $1/4$ of the expected optimal welfare when the agents' valuations are drawn from a product distribution over subadditive valuations. In addition, using techniques similar to those in the proof of Theorem 8.2, we can achieve this with uniform prices.

THEOREM 8.4. *Let $\mathcal{F} = F_1 \times \dots \times F_n$ be a product distribution over subadditive valuations (over identical items). Let OPT be the expected maximal social welfare. There exists a uniform price on the items for which the expected social welfare is at least $OPT/4$.*

Our results do not hold for general valuations, even if we can control the order of arrival.

PROPOSITION 8.5. *There is a market with a distribution over valuations (over identical items) with n agents and $m = n^2$ items such that no static pricing \mathbf{p} yields expected welfare more than $\Theta(1/n)$ of the optimal expected welfare, even if we can control the arrival order.*

PROOF. Consider a market with n identical agents (so that the arrival order does not affect the expected welfare). Each agent is unit-demand with value 1 with probability $1 - 1/n$, and single-minded with value m over the grand bundle with probability $1/n$. With probability at least $1 - 1/e$, there exists a single-minded agent, implying that the expected optimal welfare is at least $(1 - 1/e)m = \Theta(n^2)$. If we price all items at least 1 , then a single-minded agent will not buy anything, and the welfare is at most n . Otherwise, we price some item below 1 . With probability $1 - 1/n$, the first agent will buy this item and the welfare is again at most n . So the expected welfare in this case is at most $(1 - 1/n)n + (1/n)m = O(n)$. \square

9 CONCLUSION

In this article, we have studied the fraction of the optimal social welfare that can be achieved via posted prices in markets with identical items under various assumptions on the designer's information and agents' valuations. We have shown that in the Bayesian setting, uniform posted prices can guarantee $1/2$ and $1/4$ of the optimal welfare for XOS and subadditive valuations, respectively. If the designer has full information on agents' valuations, then $1/3$ of the optimal welfare can be

obtained via uniform prices for subadditive valuations. For general valuations, we exhibited a tight bound of $1/m$ for both uniform and non-uniform prices; however, if the designer can control the arrival order, then $1/2$ of the optimal welfare can be guaranteed for such valuations.

Our work sheds light on the power of uniform prices for settings with identical items. For submodular valuations in the full-information setting, there is a gap between the guarantee that can be obtained by uniform and non-uniform prices, while for XOS valuations in the Bayesian setting there is no gap. It remains open whether such a gap exists for subadditive valuations, both for the full-information and the Bayesian settings. Another interesting question is whether the constant approximation guarantee provided here for subadditive valuations over identical items holds also for subadditive valuations over heterogeneous items. This problem has been raised by Feldman et al. [2015] who provide a logarithmic (in m) bound for this setting.

Beyond the approximations themselves, our work introduces several techniques for handling valuations of identical items. For submodular valuations, given their simple structure, we can work directly with the multiset of marginal values of all agents. However, subadditive and general valuations are much trickier to handle, since the marginal values of an agent are not necessarily arranged in decreasing order. We have shown that it is nevertheless useful to consider the submodular closures of these valuations, and we developed the necessary machinery to analyze the guarantees that we can obtain using the resulting prices. More generally, our article shows that the setting of multi-unit markets encompasses a rich set of structures and questions and is therefore worthy of further exploration.

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