Introduction to ML

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Cornell CS 3110 Data Structures and Functional Programming

The ML Programming Language

- General purpose programming language designed by Robin Milner in 1970
 - Meta Language for verification
- Impure Functional Programming Language
 Eager call by value evaluation
- Static strongly typed (like Java unlike C)
 - Protect its abstraction via type checking and runtime checking
- Polymorphic Type Inference
- Dialects: OCaml, Standard ML, F#

C is not Type Safe

```
int j;
union { int i, int * p } x;
x.i = 17 ;
j = *(x.p);
```

int i, *p; i =17 p = (int *) i;

Factorial in ML

let rec fac n = if n = 0 then 1 else n * fac (n - 1)

// val fac : int -> int = <fun>

let rec fac n : int = if n = 0 then 1 else n * fac (n - 1)

```
let fac n =
    let rec ifac n acc =
        if n=0 then acc else ifac n-1, n * acc
        in ifac n, 1
```

Why Study ML?

- Functional programming will make you think differently about programming
 - Mainstream languages are all about state
 - Functional programming is all about values
- ML is "cutting edge"
 - Polymorphic Type inference
 - References
 - Module system
- Practical (small) Programming Language
- New ideas can help make you a better programmer, in any language

Plan

- Basic Programming in ML
- Type Inference for ML
- ML Modules & References

Simple Types

- Booleans
 true -: bool = true
 false -: bool = false
 if ... then ... else ...types must match
- Integers

0, 1, 2,... -: int =0, 1, ... +, *, -: int * int -> int

- Strings
 "I am a string" -: string ="I am a string"
- Floats

 $1.0, 2., 3.14159, \dots :- float = 1, 2, 3.14159$

Scope Rules

- ML enforces static nesting on identifiers
 To be explained lets
- let x = e1 in e2 \equiv (λ x.e2) e1

Tuples

4, 5, "abc" :- (int*int*string)=(4, 5, "abc")

let max1 (r1, r2) : float =
 if r1 < r2 then r2 else r1
val max1: float * float -> float = fun

let args = (3.5, 4.5) val args: float * float = (3.5, 4.5)

max1 args :- float = 4.5

let $y(x_1: t_1, x_2: t_2, ..., x_n: t_n) = e$

Pattern-Matching Tuples

let x_1 : t_1 , x_2 : t_2 , ..., x_n : $t_n = e$

let max1 (pair : float * float) : float =
 let (r1, r2) = pair in
 if r1 < r2 then r2 else r1
val max1: float * float -> float = fun

let minmax (a, b) : float * float =
 if a < b then (a, b) else (b, a)
val minmax: float * float -> float *float= fun

```
let (mn, mx) = minmax (2.0, 1.0)
val mn float 1
val mx float 2
```

The compiler guarantees the absence of runtime errors

User-Defined Types

type day = Sun | Mon | Tue | Wed | Thu | Fri | Sat

```
let int_to_day (i : int) : day =
   match i mod 7 with
    0 -> Sun
    | 1 -> Mon
    | 2 -> Tue
    | 3 -> Wed
    | 4 -> Thu
    | 5 -> Fri
    |_ -> Sat
```

Records

type person = {first:string; last:string; age:int}

```
{first="John"; last="Amstrong"; age=77}
:- person ={first="John; last="Amstrong"; age=77}
```

```
{first="John"; last="Amstrong"; age=77}.age
:- int = 77
```

```
let ja = {first="John"; last="Amstrong"; age=77}
val ja : person = {first="John"; last="Amstrong"; age=77}
```

```
let = {first=first; last=last} = ja
val first:string="John"
val last:string ="Amstrong"
```

Variant Records

• Provides a way to declare Algebraic data types

type expression = Number of int | Plus of expression * expression

```
let rec eval_exp (e : expression) : int =
  match e with
   Number(n) -> n
   | Plus (left, right) -> eval_exp(left) + eval_exp(right)
val eval_exp : expression -> int = <fun>
```

eval_exp (Plus(Plus(Number(2), Number(3)), Number(5)))
:- int = 10

Variant Records in C

```
struct exp {
    int tag ; /* Select between cases */
    union {
        struct number { int : number; }
        struct plus { struct exp *left, *right; }
        }
    }
}
```

Scope

- Local nested scopes
- Let constructs introduce a scope

let f x = e1 in e2

let x = 2and y = 3in x + y

let rec even $x = x = 0 \parallel \text{odd} (x-1)$ and odd $x = \text{not} (x = 0 \parallel \text{not} (\text{even} (x-1)))$ in odd 3110

Polymorphism

- A Polymorphic expression may have many types
- There is a "most general type"
- The compiler infers types automatically
- Programmers can restrict the types
- Pros:
 - Code reuse
 - Guarantee consistency
- Cons:
 - Compile-time
 - Some limits on programming

```
let max1 (r1, r2) =
  if r1 < r2 then r2 else r1
val max1: 'a * 'a -> 'a = fun
```

max1 (5, 7)

$$: - int = 7$$

max1 (5, 7.5)

Polymorphic Lists

[] - : 'a list = []

[2; 7; 8] - : int list = [2; 7; 8]

[(2, 7) ; (4, 9) ; 5] Error: This expression has type int but an expression was expected of type int * int

Functions on Lists

```
let rec length 1 =
    match 1 with
    [] -> 0
    | hd :: t1 -> 1 + length t1
val length : 'a list -> int = <fun>
```

```
length [1; 2; 3] + length ["red"; "yellow"; "green"]
:- int = 6
```

length ["red"; "yellow"; 3]

Higher Order Functions

- Functions are first class objects
 - Passed as parameters
 - Returned as results
- Practical examples
 - Google map/reduce

Map Function on Lists

• Apply function to every element of list

```
let rec map f arg = function
   [] -> []
   | hd :: tl -> f hd :: (map f tl)
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>
```

map (fun x -> x+1)
$$[1;2;3]$$
 \longrightarrow $[2,3,4]$

• Compare to Lisp

```
(define map
    (lambda (f xs)
        (if (eq? xs ()) ()
            (cons (f (car xs)) (map f (cdr xs)))
    )))
```

More Functions on Lists

Append lists

let rec append l1 l2 =
 match l1 with
 [] -> l2
 [hd :: tl -> hd :: append (tl l2)

• Reverse a list

let rec reverse l = function
 [] -> []
 [hd :: tl -> append (reverse tl) [hd]

• Questions

– How efficient is reverse?

– Can it be done with only one pass through list?

More Efficient Reverse





Currying

let plus (x, y) = x + yval plus : int * int -> int = fun

let plus (z : int * int) = match z with $(x, y) \rightarrow x + y$

let plus = fun (z : int * int) -> match z with (x, y) -> x + y

let plus x y = x + yval plus : int -> int -> int

let p1 = plus 5val p1 : int -> int = fun

let p2 = p1 7 val p2 : int = 12

Functional Programming Languages

PL	types	evaluation	Side-effect
scheme	Weakly typed	Eager	yes
ML OCAML F#	Polymorphic strongly typed	Eager	References
Haskell	Polymorphic strongly typed	Lazy	None

Things to Notice

• Pure functions are easy to test

prop_RevRev l = reverse(reverse l) == l

- In an imperative or OO language, you have to
 - set up the state of the object and the external state it reads or writes
 - make the call
 - inspect the state of the object and the external state
 - perhaps copy part of the object or global state, so that you can use it in the post condition

Things to Notice

Types are everywhere.

reverse:: [w] -> [w]

- Usual static-typing panegyric omitted...
- In ML, types express high-level design, in the same way that UML diagrams do, with the advantage that the type signatures are machine-checked
- Types are (almost always) optional: type inference fills them in if you leave them out

Recommended ML Textbooks

- L. C. PAULSON: ML for the Working Programmer
- J. Ullman: Elements of ML Programming
- R. Harper: Programming in Standard ML

Recommended Ocaml Textbooks

• Xavier Leroy: The OCaml system release 4.02

– Part I: Introduction

- Jason Hickey: Introduction to Objective Caml
- Yaron Minsky, Anil Madhavapeddy, Jason Hickey: Real World Ocaml

Summary

- Functional programs provide concise coding
- Compiled code compares with C code
- Successfully used in some commercial applications
 - F#, ERLANG, Jane Street
- Ideas used in imperative programs
- Good conceptual tool
- Less popular than imperative programs